

GOTO Copenhagen 2018 Conference Nov. 19 - 21

Turning the JVM into a Polyglot VM with Graal





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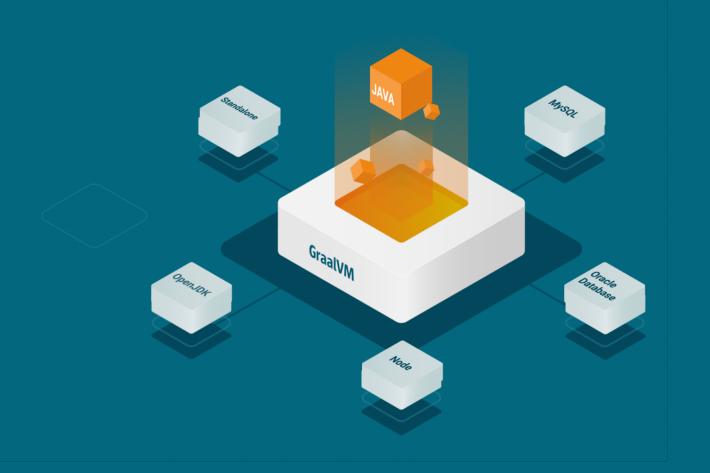


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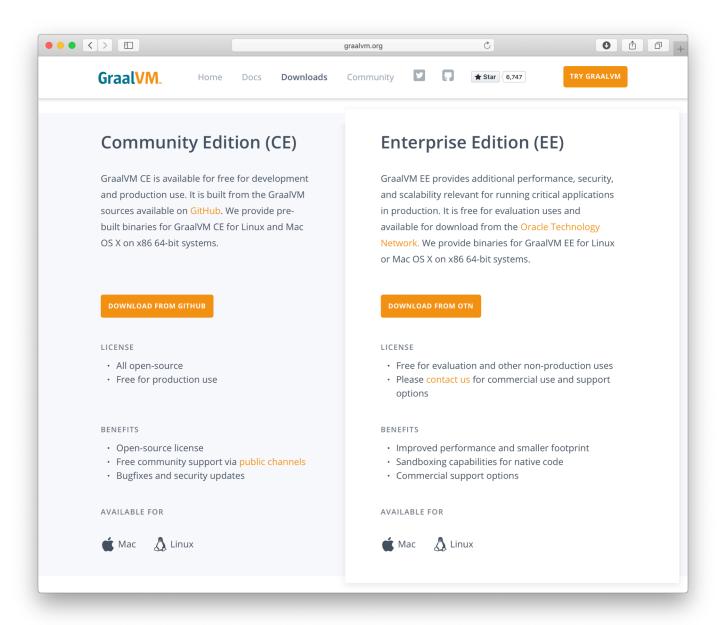
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GraalVM...

Run Programs Faster Anywhere



- What, in concrete terms, is GraalVM?
- What can I practically do with it?
- What do all these things have to do with each other?
- What is the big idea?
- What kind of change is this going to enable?





Using GraalVM as your JDK

- You can use GraalVM as a drop-in replacement for OpenJDK
- 1.8 at the moment, will be updated to the 11 LTS soon
- Includes all the same commands, flags, options and so on

Add it to your \$PATH

\$ export PATH=graalvm-ee-1.0.0-rc8/Contents/Home/bin:\$PATH



```
public static void main(String[] args) {
   Arrays.stream(args)
            .flatMap(TopTen::fileLines)
            .flatMap(line -> Arrays.stream(line.split("\\b")))
            .map(word -> word.replaceAll("[^a-zA-Z]", ""))
            .filter(word -> word.length() > 0)
            .map(word -> word.toLowerCase())
            .collect(Collectors.groupingBy(Function.identity(), Collectors.counting()))
            .entrySet().stream()
            .sorted((a, b) -> -a.getValue().compareTo(b.getValue()))
            .limit(10)
            .forEach(e -> System.out.format("%s = %d%n", e.getKey(), e.getValue()));
```

Compile and run as normal

```
$ javac TopTen.java
$ time java TopTen large.txt
...
real 0m18.905s

This demo is
run with the
```



EE version

Compare to standard OpenJDK

```
$ time java -XX:-UseJVMCICompiler TopTen large.txt
```

•••

real 0m23.102s



What is going on?

- The Graal just-in-time compiler is one part of GraalVM
- It replaces (or runs as a tier above) the existing JIT compilers like C2
- It's written in Java, which we think lets us improve it more easily, so it achieves better performance than C2
- Here we're getting 20% faster performance on a benchmark
- Twitter see 18% faster in production on real Scala applications, using only the CE version EE not needed for high performance
- Is it odd that a JIT compiler for Java is written in Java?



```
synchronized (oject) {
    . . . A
synchronized (oject) {
    . . . B
synchronized (oject) {
    ...B
```

```
for (MonitorExitNode monitorExitNode: graph.getNodes(MonitorExitNode.TYPE)) {
    FixedNode next = monitorExitNode.next():
    if ((next instanceof MonitorEnterNode || next instanceof RawMonitorEnterNode)) {
        // should never happen, osr monitor enters are always direct successors of the graph
       // start
        assert !(next instanceof OSRMonitorEnterNode);
        AccessMonitorNode monitorEnterNode = (AccessMonitorNode) next;
        if (isCompatibleLock(monitorEnterNode, monitorExitNode)) {
            /*
             * We've coarsened the lock so use the same monitor id for the whole region,
            * otherwise the monitor operations appear to be unrelated.
             */
           MonitorIdNode enterId = monitorEnterNode.getMonitorId();
           MonitorIdNode exitId = monitorExitNode.getMonitorId();
            if (enterId != exitId) {
                enterId.replaceAndDelete(exitId);
           GraphUtil.removeFixedWithUnusedInputs(monitorEnterNode);
           GraphUtil.removeFixedWithUnusedInputs(monitorExitNode);
```

JVMCI is the interface that lets you plug in a new JIT

```
$ time java -XX:-UseJVMCICompiler TopTen large.txt
...
real 0m23.102s
```



Graal also works on standard OpenJDK 11

- Graal (the JIT compiler part) is also included in Graal
- As an experimental, unsupported option, hidden behind flags
- Older version, due to the release cycle
- We'd recommend using the GraalVM package to experiment with

Enable Graal in OpenJDK 11

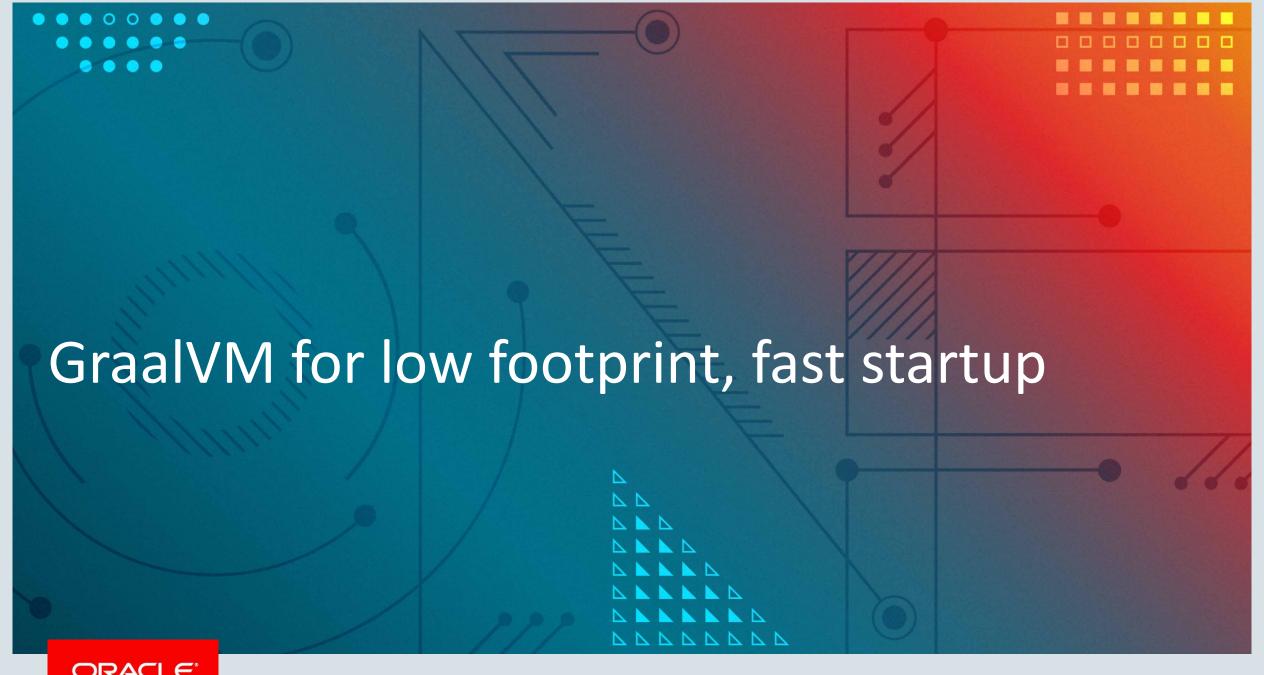
On AMD64, on macOS and Linux

```
$ java -XX:+UnlockExperimentalVMOptions \
-XX:+EnableJVMCI \
-XX:+UseJVMCICompiler \
```

• • •

Doesn't mention Graal, does it? JVMCI does service discovery and automatically finds Graal as the only JVMCI compiler available





What about shorter running applications or functions?

- The JVM typically has a relatively slow time to start
 - Compared to simpler VMs, like Python or Ruby
 - Compared to native executables like those produced from Go or Rust
 - JRuby 'hello world' startup time is an order of magnitude worse than standard Ruby
- The JVM typically takes up a relatively large amount of disk space
 - Can be helped with jlink down to tens of MB
- The JVM typically takes up a relatively large amount of RAM
 - Interpreter, compiler, classfile parser, verifier etc all take up space



Run as normal

```
$ time java TopTen small.txt
...
real 0m0.408s
```

Compile to native using GraalVM

```
$ native-image TopTen
...
$ time ./topten small.txt
...
real 0m0.112s
```



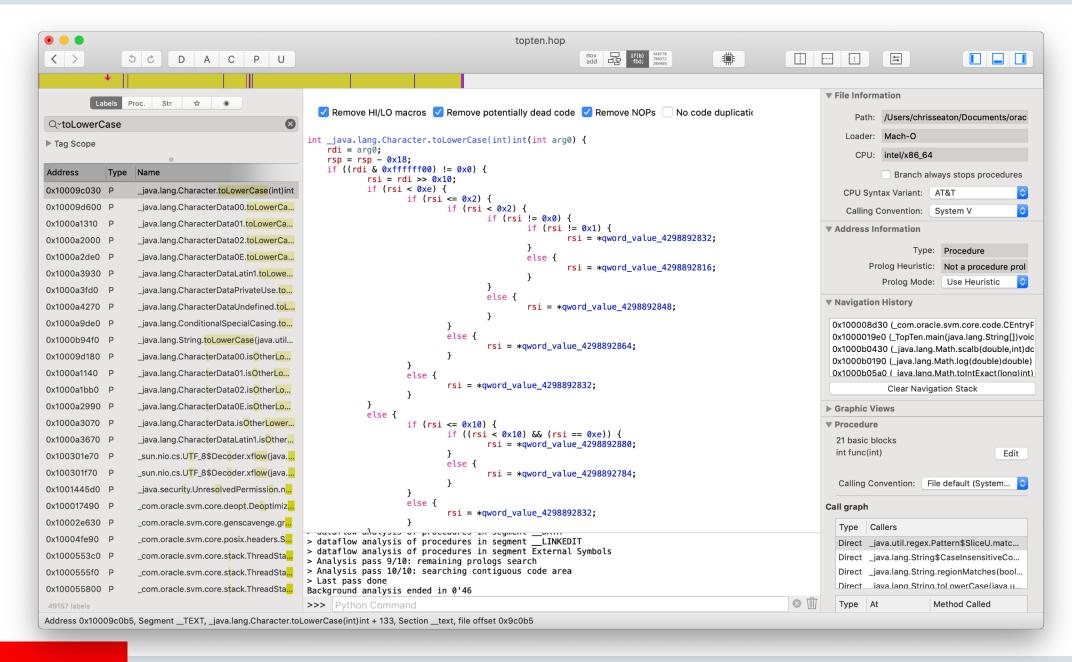




\$ du -h topten
8.8M topten



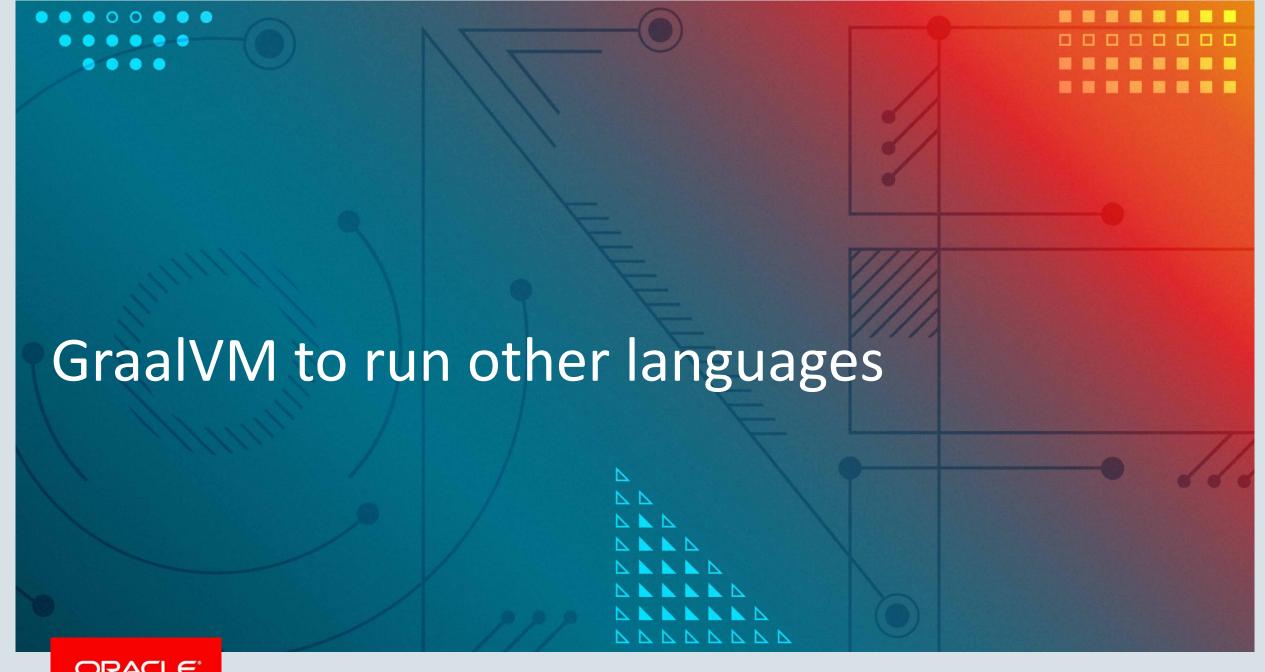
```
$ otool -L topten
topten:
/usr/lib/libSystem.B.dylib
/usr/lib/libz.1.dylib
/System/Library/CoreFoundation
```



So what does this have to do with a JIT?

- Graal is written in Java
- So it can be used as a library from other Java code
- We realized that we could write a program to use it ahead-of-time, to build and ahead-of-time compiler
- This is what the native-image tool is a Java application that uses Graal as a library

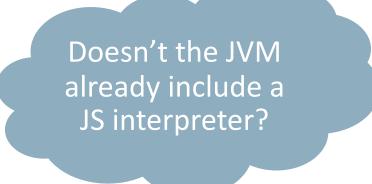




GraalVM includes a new JavaScript interpreter

```
$ js -version
Graal JavaScript 1.0 (GraalVM CE Native 1.0.0-rc8)

$ js
> print("hello");
hello
Doesn't the JVM
```



GraalVM also includes an implementation of Node.js

```
$ node --version
v10.9.0
```

```
$ npm --version
6.2.0
```

```
var express = require('express');
var app = express();
app.get('/', function (req, res) {
  res.send('<h1>Hello!</h1>');
});
app.listen(8080, function () {
  console.log('serving at http://localhost:8080')
});
```

```
$ npm install express
$ node hello-express.js
serving at http://localhost:8080
```

You can plug new languages into GraalVM

```
$ gu install ruby
$ gu install python
$ gu install R
```

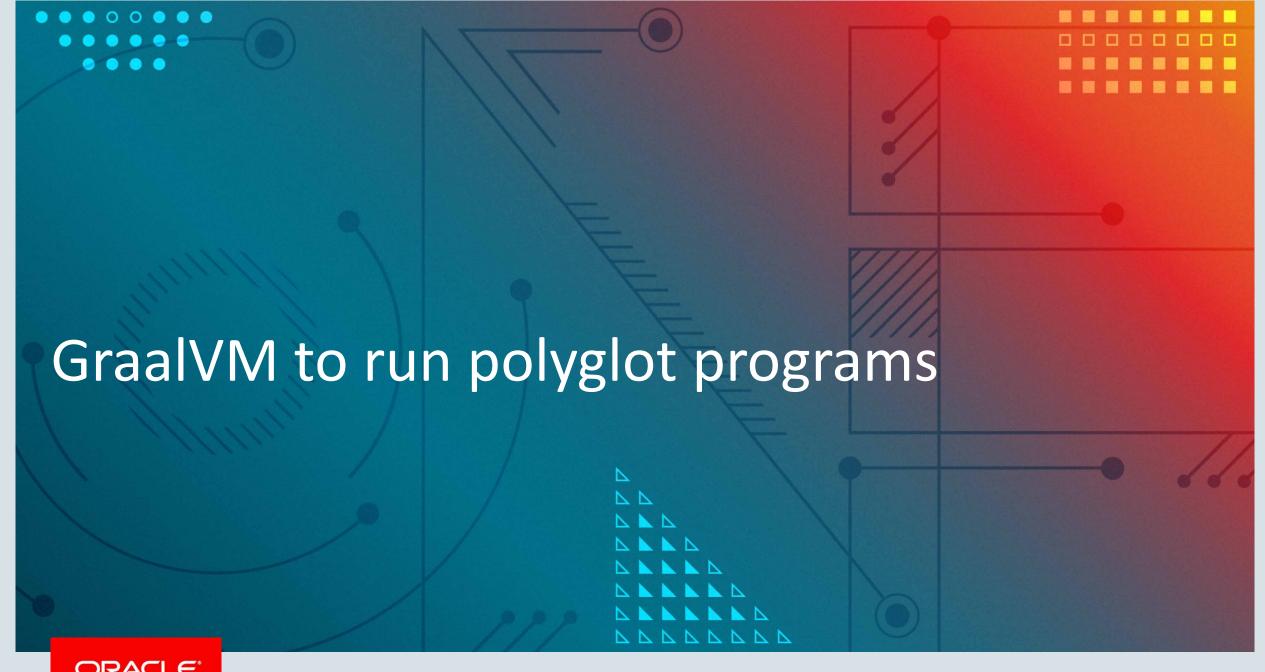
You can plug new languages into GraalVM

```
$ ruby --version
truffleruby 1.0.0-rc8, like ruby 2.4.4, GraalVM CE Native [x86_64-darwin]
$ graalpython --version
Graal Python 3.7.0 (GraalVM CE Native 1.0.0-rc8)
$ R --version
R version 3.4.0 (FastR)
```

So what does this have to do with a JIT?

- We realized instead of writing languages that emit bytecode at runtime (JRuby) we could write languages that use the Graal JIT directly
- But that's hard, so we realized we could write a framework, Truffle, to do that automatically, based on a simple interpreter
- Faster, as they use a more powerful JIT more directly
- Simpler, because a framework does most of the hard work, so easy to implement lots of languages
- Interopable (polyglot) because they all use the same system
- Using native-image they start quickly





GraalVM is polyglot as well as multi-language

- Like many languages implemented on the JVM, our languages can use Java libraries
- Run ruby-java.rb

```
BigInteger = Java.type('java.math.BigInteger')
puts BigInteger.valueOf(2).pow(100).toString
```

You need to run ruby with the --jvm flag



GraalVM is polyglot as well as multi-language

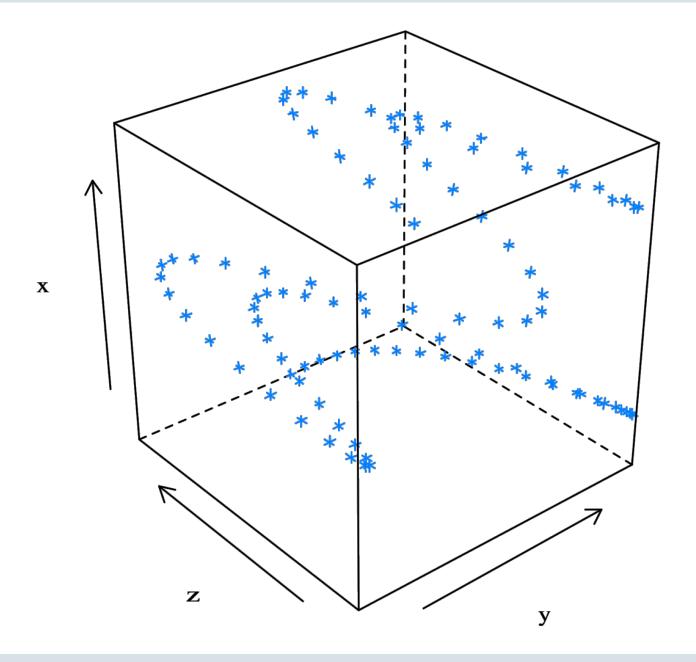
• Run python-java.rb

```
import java
BigInteger = java.type('java.math.BigInteger')
print(BigInteger.valueOf(2).pow(100).toString())
```

You need to run graalpython with the --jvm flag

Java is just another language in this case

```
const express = require('express');
const app = express();
                                               You need to run node
app.get('/', function (req, res) {
                                                with the --jvm flag
  res.send(Interop.eval('R',
    `svg();
     require(lattice);
    x < -1:100
    y \ll \sin(x/10)
     z \leftarrow cos(x^1.3/(runif(1)*5+10))
    print(cloud(x~y*z, main="cloud plot"))
    grDevices:::svg.off()
    `));
})
app.listen(8080, function () {
  console.log('serving at http://localhost:8080');
});
```



So what does this have to do with a JIT?

- All the languages use the same JIT, and the same high-level implementation framework, so they can all work together
- Integration is at a higher level than with traditional bytecode implementation



Is our approach only suited to Java and dynamic languages?

- There's nothing special about native languages
- C has the same if statements and while loops Ruby does
- C has pointers and malloc, but so does Ruby in its FFI module

Is our approach only suited to Java and dynamic languages?

- Example running gzip on the JVM
 - Not a clean piece of code
 - -8.6 k lines of C
 - Macros, pointer arithmetic, unions
 - We'll avoid the complexity of autotools and make by using a single-file version
 - http://people.csail.mit.edu/smcc/projects/single-file-programs/



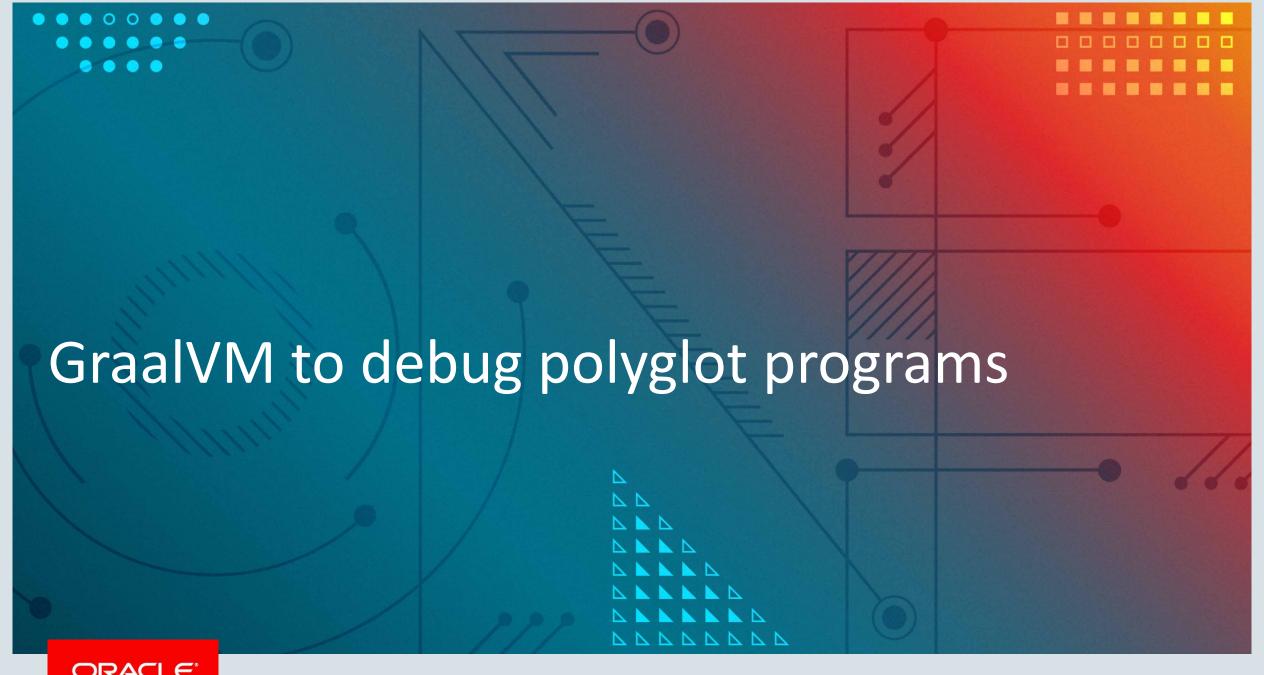
```
if (more == (unsigned)EOF) {
    /* Very unlikely, but possible on 16 bit machine if strstart == 0
     * and lookahead == 1 (input done one byte at time)
     */
    more--;
} else if (strstart >= WSIZE+MAX_DIST) {
    /* By the IN assertion, the window is not empty so we can't confuse
     * more == 0 with more == 64K on a 16 bit machine.
     */
    Assert(window_size == (ulg)2*WSIZE, "no sliding with BIG_MEM");
    memcpy((char*)window, (char*)window+WSIZE, (unsigned)WSIZE);
    match_start -= WSIZE;
    strstart -= WSIZE; /* we now have strstart >= MAX_DIST: */
   if (rsync_chunk_end != 0xFFFFFFFUL)
       rsync_chunk_end -= WSIZE;
    block_start -= (long) WSIZE;
    for (n = 0; n < HASH_SIZE; n++) {</pre>
        m = head[n];
        head[n] = (Pos)(m >= WSIZE ? m-WSIZE : NIL);
    for (n = 0; n < WSIZE; n++) {</pre>
        m = prev[n];
        prev[n] = (Pos)(m >= WSIZE ? m-WSIZE : NIL);
        /* If n is not on any hash chain, prev[n] is garbage but
         * its value will never be used.
         */
    more += WSIZE;
```

```
$ clang -c -emit-llvm gzip.c
$ gzip small.txt
$ lli gzip.bc -d small.txt.gz
```

So what does this have to do with a JIT?

- We can use the JIT that we use for Java, JavaScript, Ruby, Python, R and so on, for C as well
- Actually any language that can target LLVM
- C, C++, Objective C, Swift, Fortran, Rust, etc
- Genuine potential for dynamic optimization
- Potential for sandboxing as well

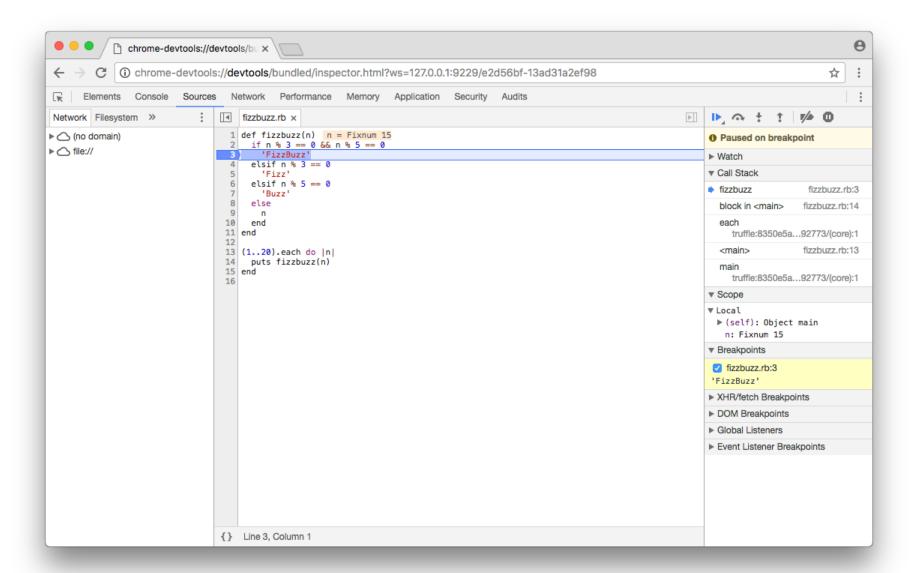


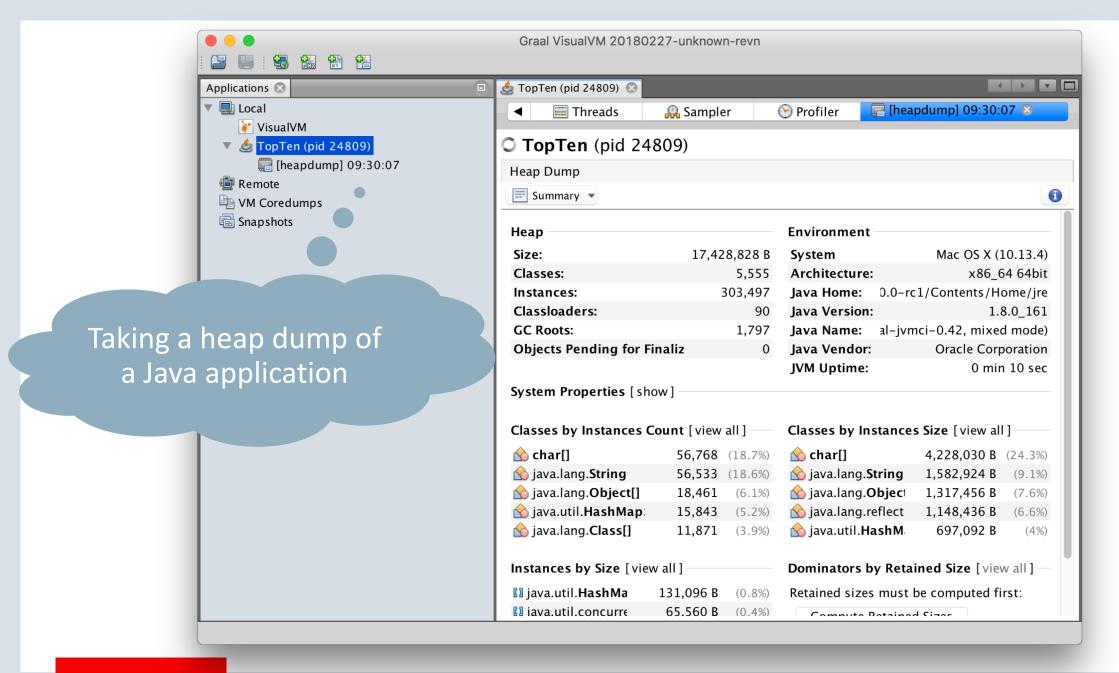


Tooling for these extra languages

- With other languages on the JVM you usually have to use a Java debugger
- Perhaps with source information in the bytecode for the guest language
- Some custom debuggers, but not for all languages

```
def fizzbuzz(n)
  if n % 3 == 0 && n % 5 == 0
    'FizzBuzz'
 elsif n % 3 == 0
    'Fizz'
 elsif n % 5 == 0
    'Buzz'
 else
    n
 end
end
(1..20).each do |n|
 puts fizzbuzz(n)
end
```

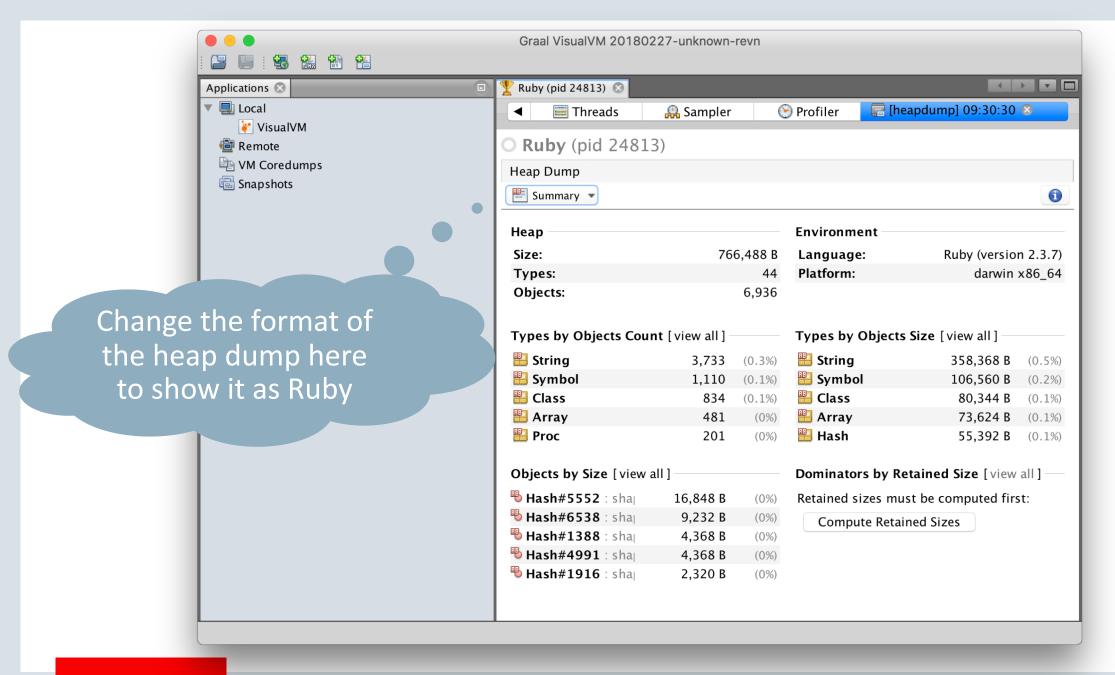




Monitoring for these extra languages

- VisualVM and other similar tools let you monitor the JVM
- Non-JVM languages often don't have this kind of tool
- With other languages on the JVM, then often show the underlying Java objects, rather than the guest language objects

- \$ ruby render.rb
- \$ jvisualvm



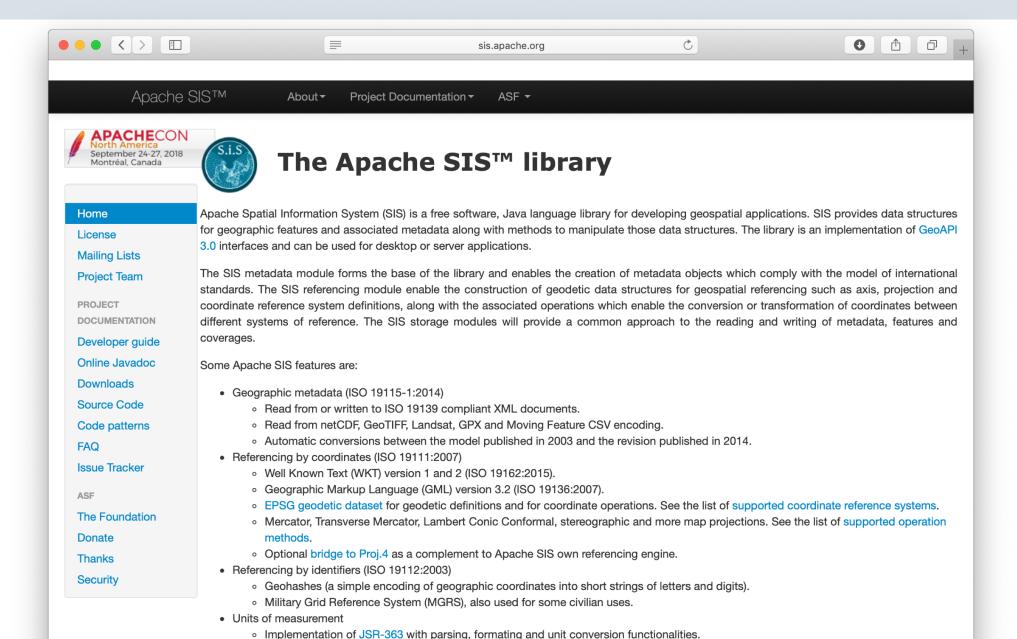
So what does this have to do with a JIT?

- All the languages are implemented in the same framework, so the debugger can understand them all via that framework
- The Graal JIT has support for deoptimization, so can debug optimized code running in production



Java code as a native library

- The Java ecosystem is phenomenal
- Often more and better libraries than available in other languages
- In the examples so far, it's always been the Java code that has owned the process
- Can we run Java code inside another application that we already have?





```
import org.apache.sis.distance.DistanceUtils;
public class Distance {
    public static void main(String[] args) {
        final double aLat = Double.parseDouble(args[0]);
        final double aLong = Double.parseDouble(args[1]);
        final double bLat = Double.parseDouble(args[2]);
        final double bLong = Double.parseDouble(args[3]);
        System.out.printf("%f km%n", DistanceUtils.getHaversineDistance(aLat, aLong, bLat, bLong));
```

```
$ native-image -cp sis.jar:. Distance
...
$ ./distance 51.507222 -0.1275 40.7127 -74.0059
5570.25 km
```

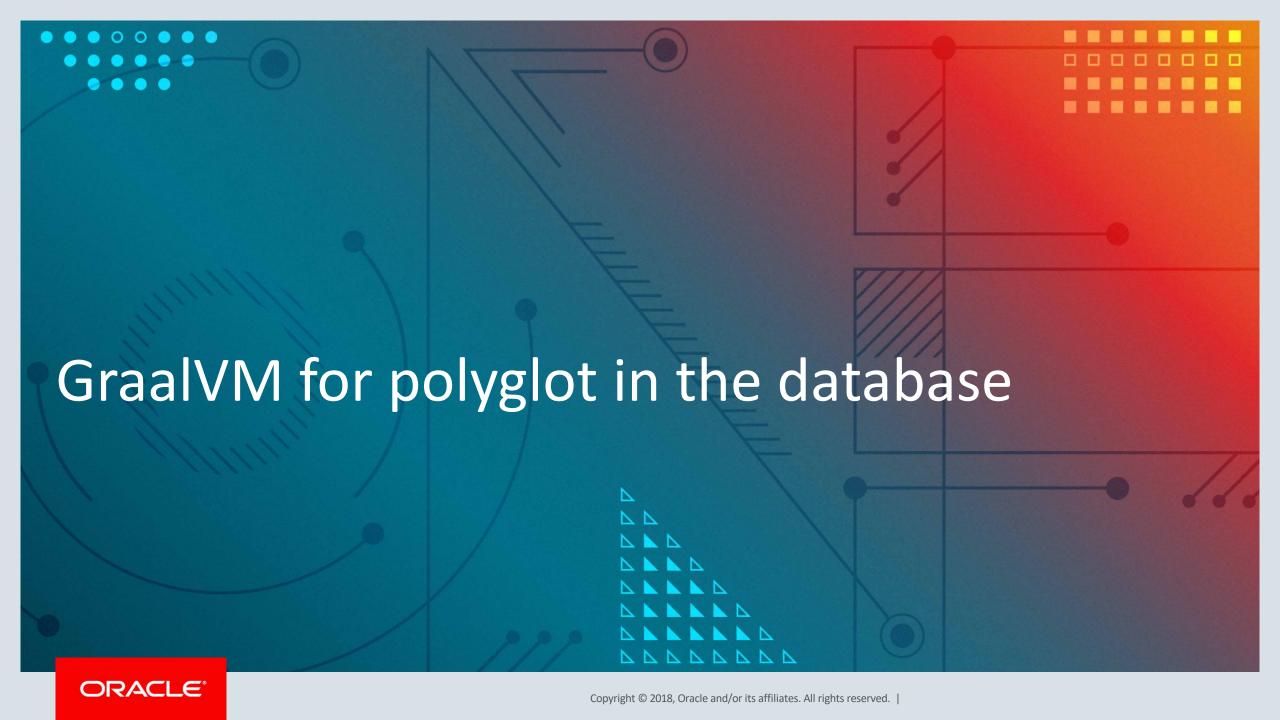
```
. . .
import org.graalvm.nativeimage.IsolateThread;
import org.graalvm.nativeimage.c.function.CEntryPoint;
public class Distance {
    . . .
    @CEntryPoint(name = "distance")
    public static double distance(IsolateThread thread,
          double a_lat, double a_long,
          double b_lat, double b_long) {
        return DistanceUtils.getHaversineDistance(a_lat, a_long, b_lat, b_long);
```

\$ native-image -cp sis.jar:. -H:Kind=SHARED_LIBRARY \
 -H:Name=libdistance

```
#include <stdlib.h>
#include <stdio.h>
#include <libdistance.h>
int main(int argc, char **argv) {
  graal_isolate_t *isolate = NULL;
  graal_isolatethread_t *thread = NULL;
  if (graal_create_isolate(NULL, &isolate) != 0 || (thread = graal_current_thread(isolate)) == NULL) {
    fprintf(stderr, "initialization error\n");
    return 1;
  double a_lat = strtod(argv[1], NULL);
  double a_long = strtod(argv[2], NULL);
  double b_lat = strtod(argv[3], NULL);
  double b_long = strtod(argv[4], NULL);
  printf("%f km\n", distance(thread, a_lat, a_long, b_lat, b_long));
  return 0;
```

```
$ clang -I. -L. -ldistance distance.c -o distance
$ otool -L distance
distance:
    libdistance.dylib
    /usr/lib/libSystem.B.dylib
$ ./distance 51.507222 -0.1275 40.7127 -74.0059
5570.25 km
```

libdistance.dylib



Demo using the Oracle Database MLE

- Multi-lingual (polyglot) edition
- Available as a Docker image
- Subject to the Oracle Technology Network license agreement, so you need to accept that and download it yourself

https://oracle.github.io/oracle-db-mle/releases/0.2.7/docker/

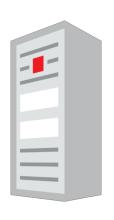


```
$ docker load --input mle-docker-0.2.7.tar.gz # takes a while
$ docker run mle-docker-0.2.7 # takes a while
$ docker ps
$ docker exec -ti <container_id> bash -li
```

JavaScript in the client and frontend, Oracle in the backend





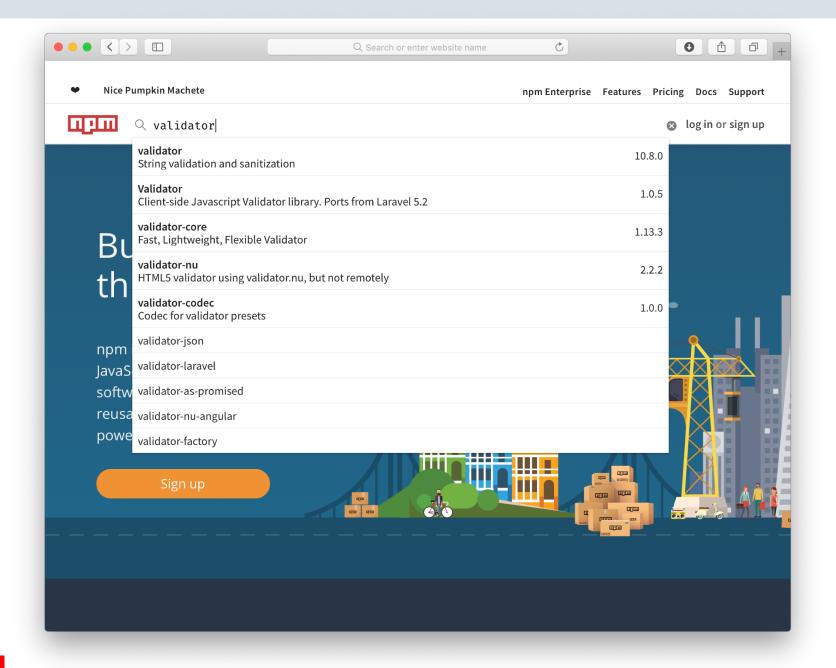


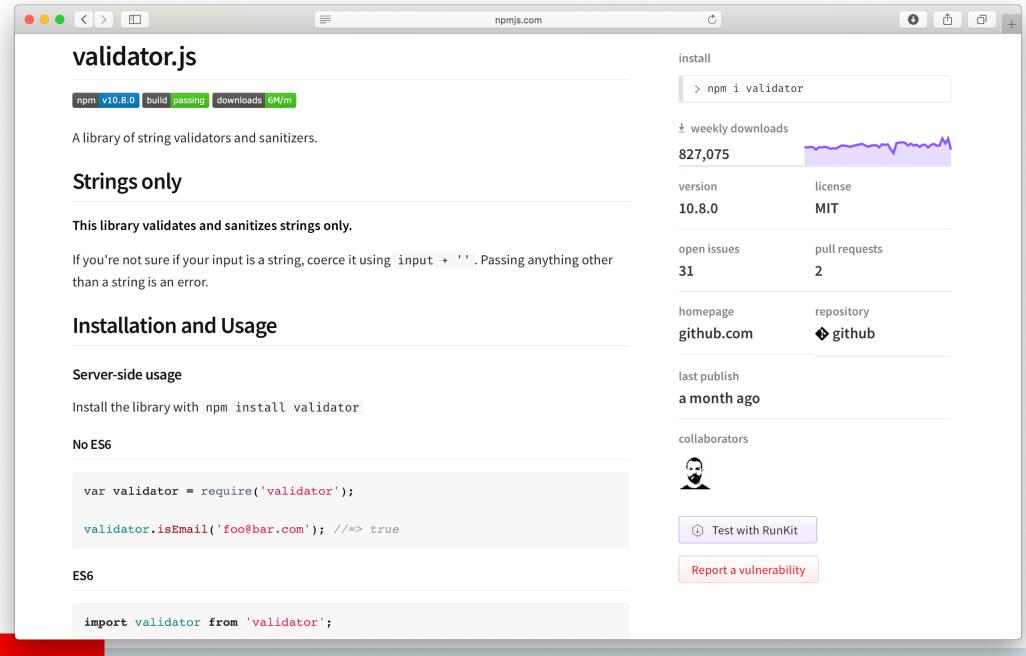












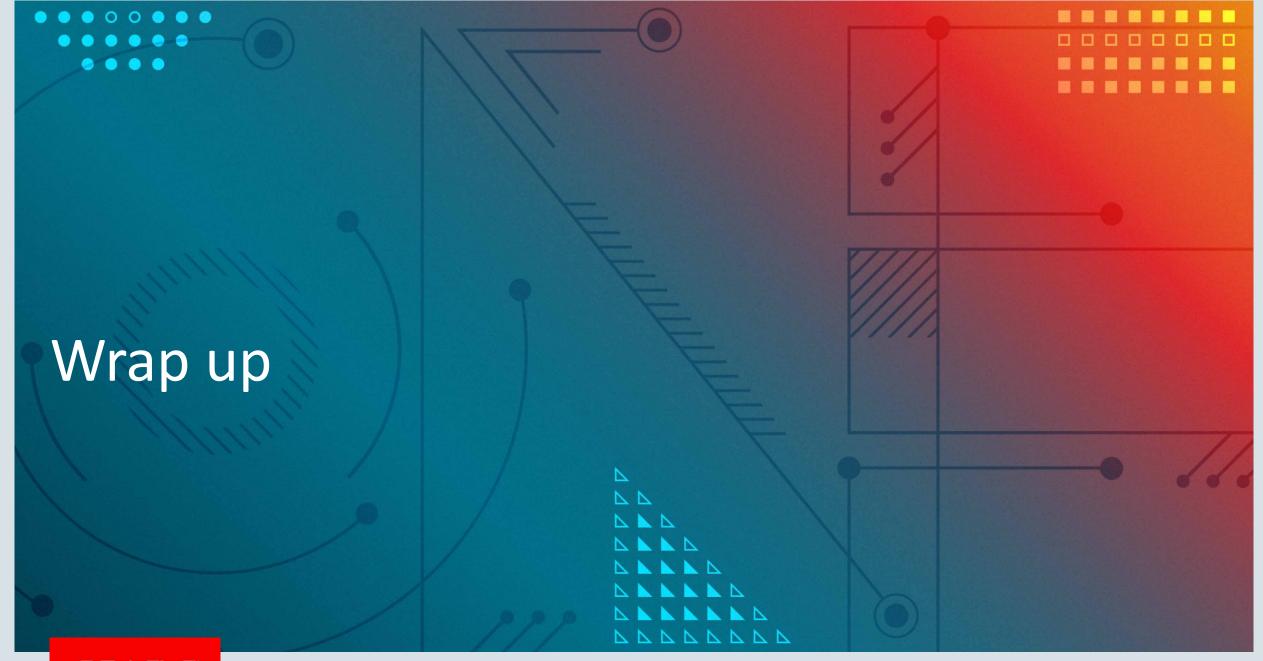
```
$ echo "{}" > package.json
$ npm install validator
$ npm install @types/validator
$ dbjs deploy -u scott -p tiger -c localhost:1521/ORCLCDB validator
$ sqlplus scott/tiger@localhost:1521/ORCLCDB
```

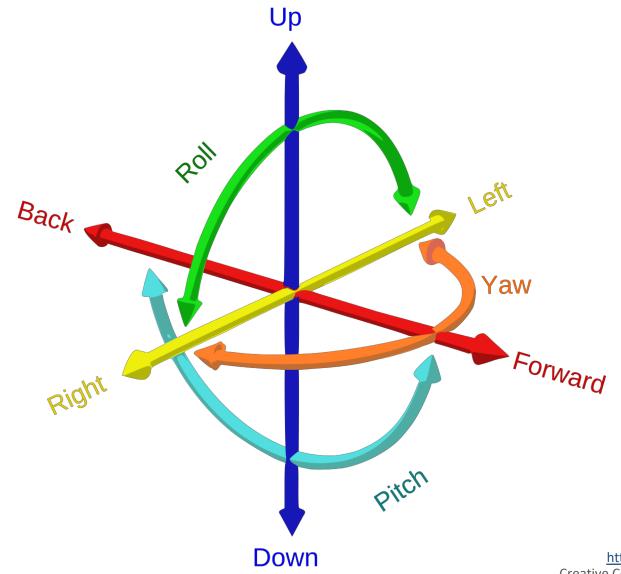
```
SQL> select validator.isEmail('oleg.selaev@oracle.com') from dual;
SQL> select validator.isEmail('oleg.selaev') from dual;
```

How does this bring it all together?

- This is a...
 - JavaScript interpreter, implemented in Java, using our framework
 - Using the polyglot interface to talk to the query language
 - Compiled ahead-of-time using Graal into a native library
 - Which can be linked into the database
 - Including Graal as a JIT within that library for high performance







I think it's about giving people degrees of freedom

- Let people run the language they want
- With the ecosystem of libraries they want
- On the JVM or on native
- Embedded or embedding
- With the tooling they want
- With the performance they want
- 'One compiler to rule them all'

Get in touch with us

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