

Cloud Spanner

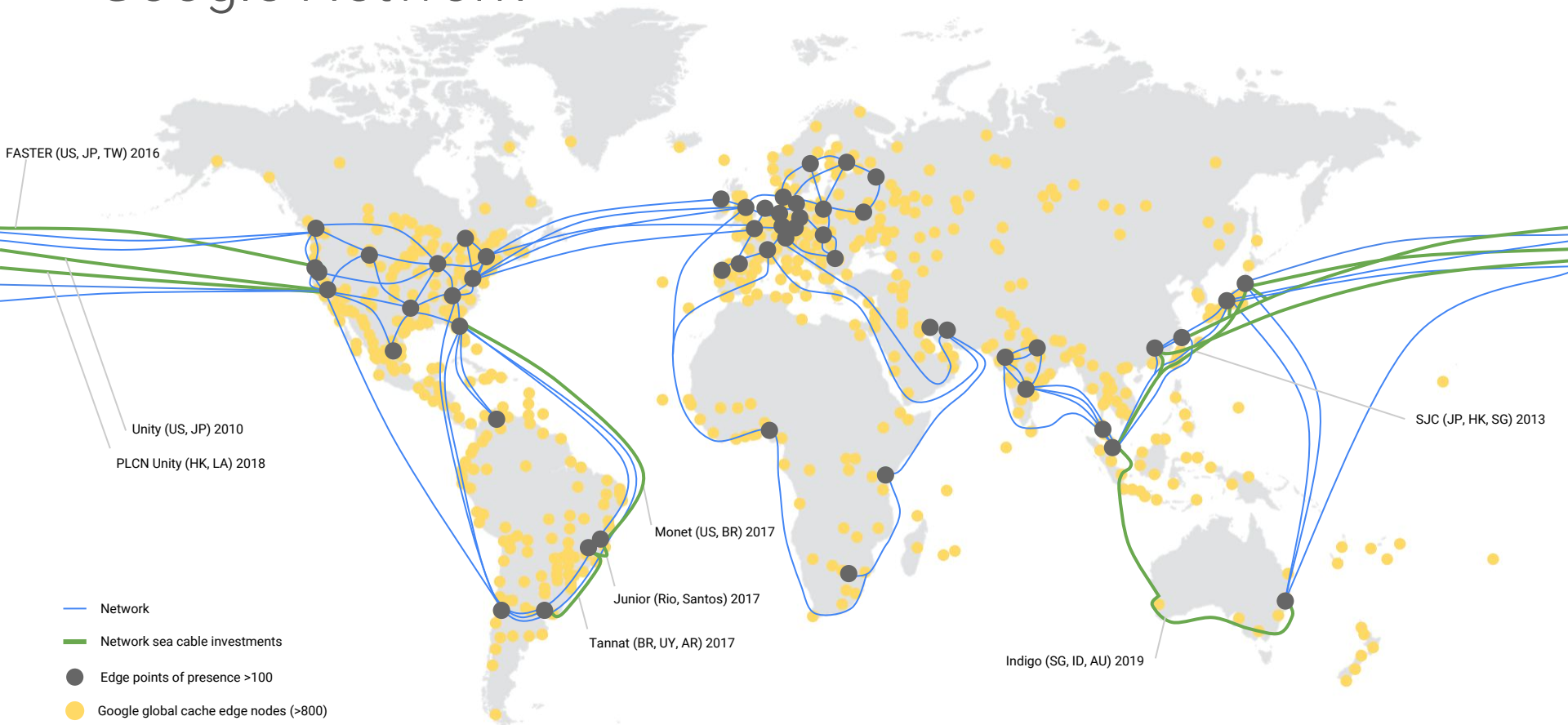
Joe Intrakamhang
Solutions Engineer
[@joe_int](#)

A wide-angle, high-angle shot of a vast data center. The ceiling is a complex network of dark metal trusses and pipes, with some fluorescent lights visible. The floor is a light-colored tiled grid. In the foreground and middle ground, there are numerous rows of server racks. Some racks are illuminated with bright blue light, while others are dimly lit. The overall atmosphere is industrial and high-tech.

\$29.4 Billion

3 Year Trailing CAPEX Investment

Google Network



Background

Why build Spanner?



It's 2005...

Google's needs



Horizontally Scaling Database



ACID Transactions with global consistency



No downtime!

Existing options



Nonrelational



Relational

Overview

What is Cloud Spanner?

What is Cloud Spanner?



Google's mission-critical scalable relational Database Service



Fully managed, database service with global scale



Traditional relational semantics: schemas, ACID transactions, SQL



Automatic, synchronous replication within and across regions for availability



Battle-tested within Google for 5+ yrs (AdWords, GooglePlay)

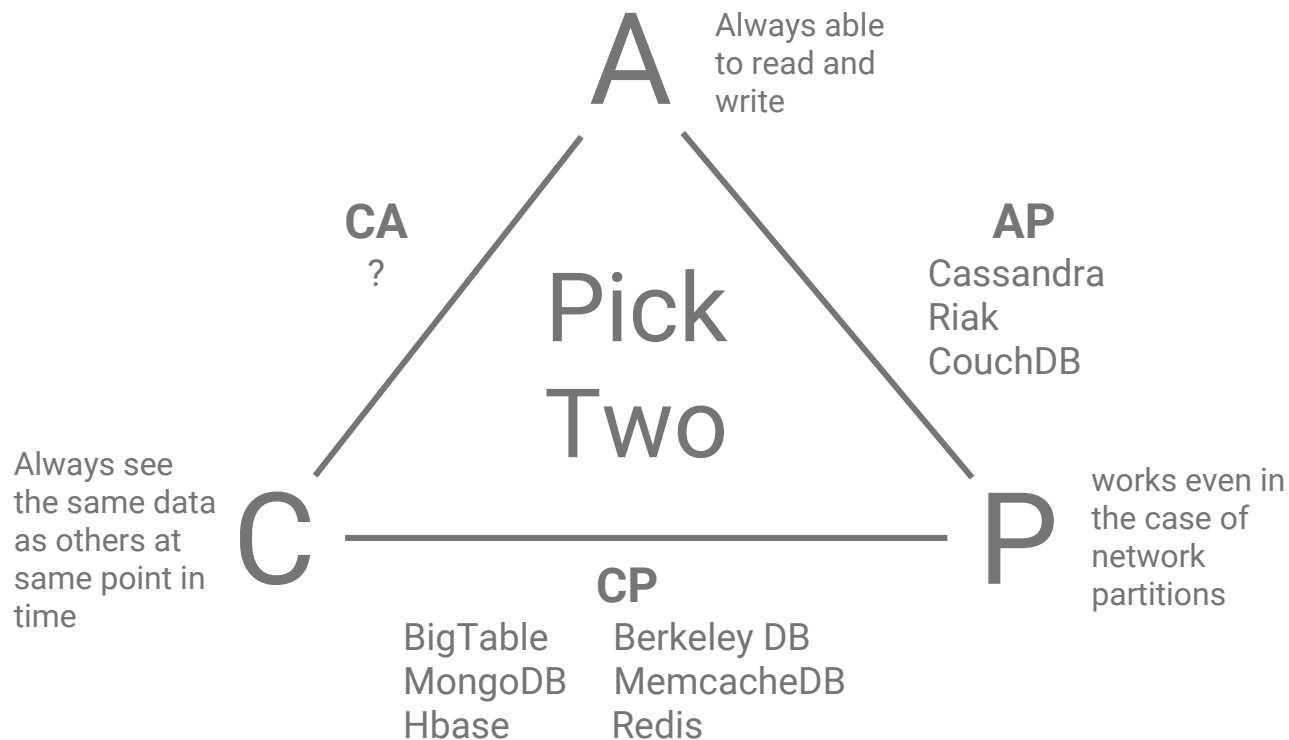


How does it compare?

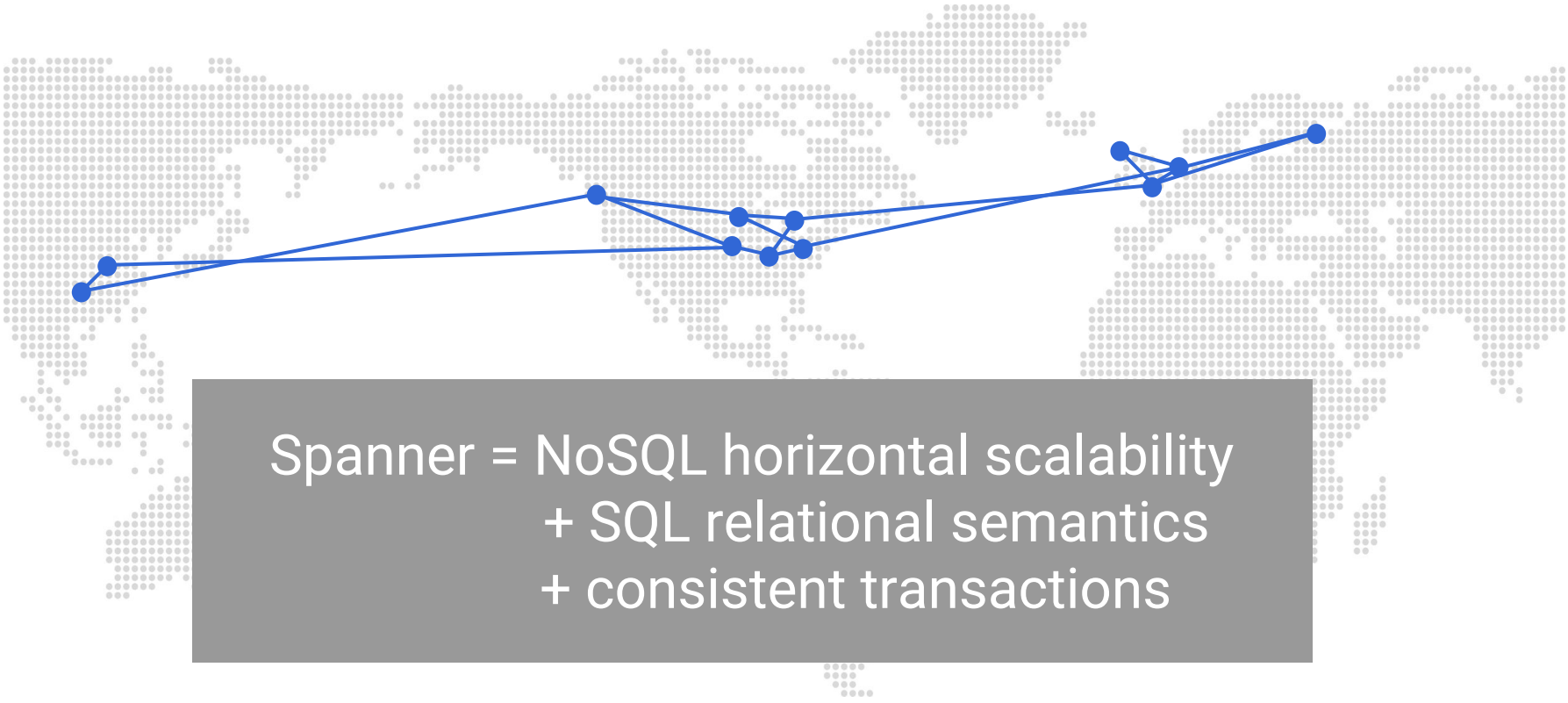


	CLOUD SPANNER	TRADITIONAL RELATIONAL	TRADITIONAL NON-RELATIONAL
Schema	✓ Yes	✓ Yes	✗ No
SQL	✓ Yes	✓ Yes	✗ No
Consistency	✓ Strong	✓ Strong	✗ Eventual
Availability	✓ High	✗ Failover	✓ High
Scalability	✓ Horizontal	✗ Vertical	✓ Horizontal
Replication	✓ Automatic	🔄 Configurable	🔄 Configurable

What alternatives are there?



Spanner tries to be both, SQL and distributed, how?



Spanner = NoSQL horizontal scalability
+ SQL relational semantics
+ consistent transactions

Open standards



Standard SQL (ANSI 2011)



Encryption, Audit logging, Identity
and Access Management



Client libraries in popular languages
(Java, C#, Python, Go, Node.js, etc.)

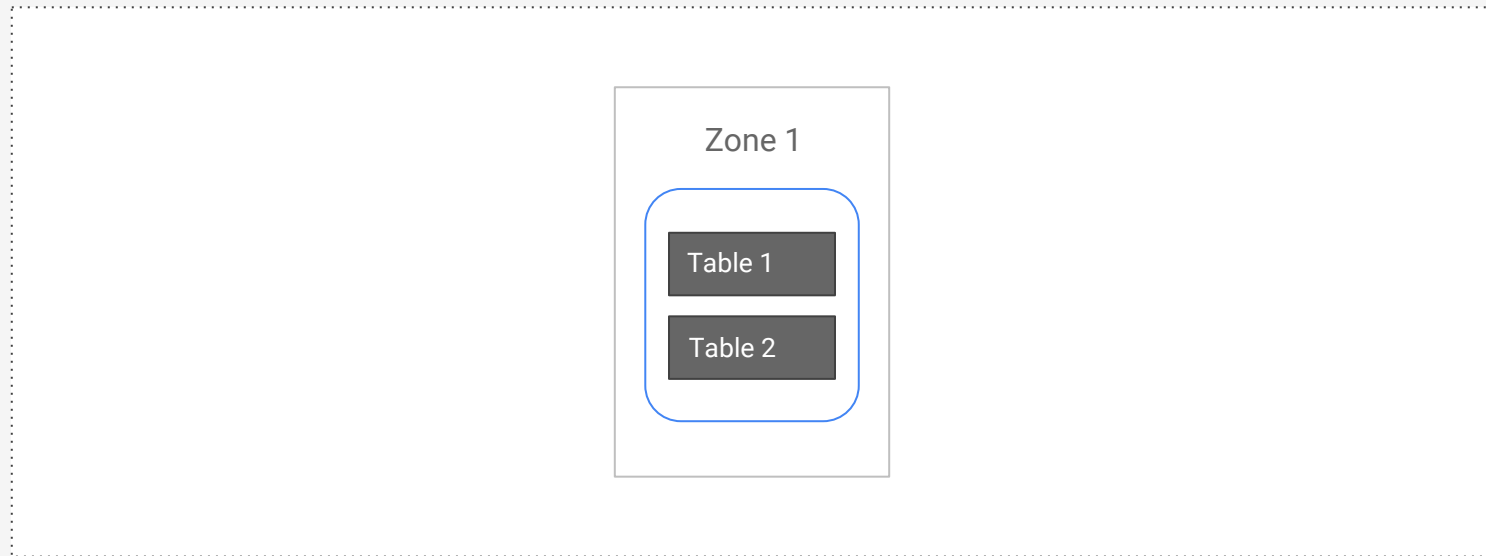


JDBC driver

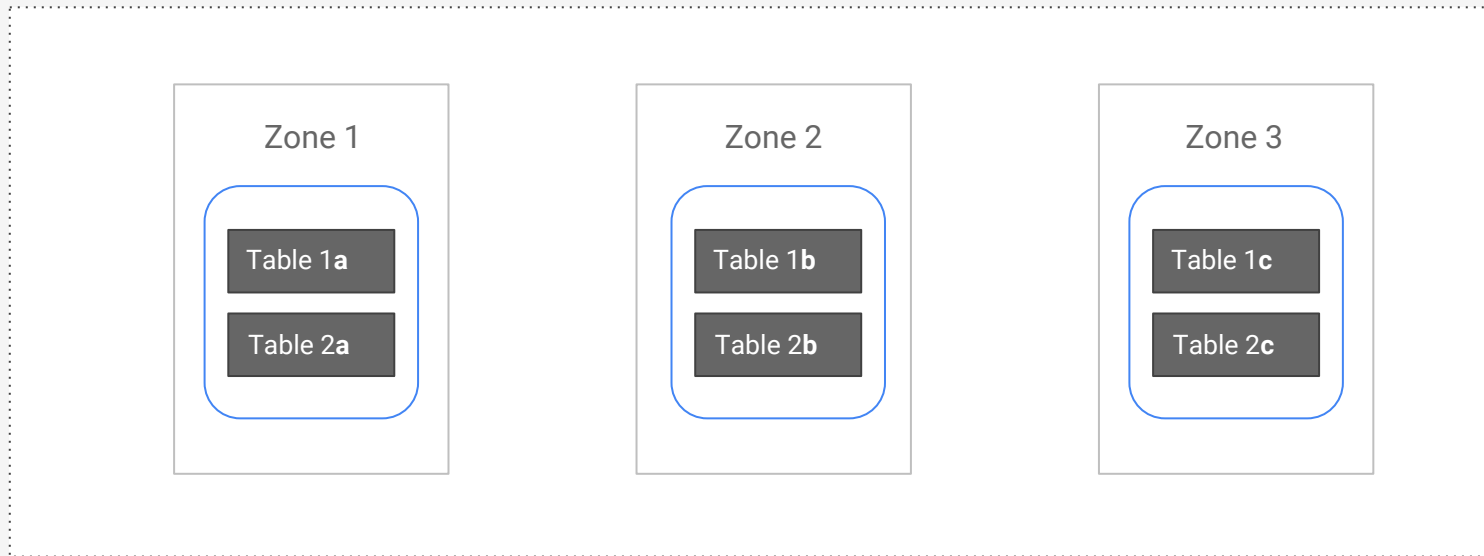


How Spanner works

Traditional Database layout



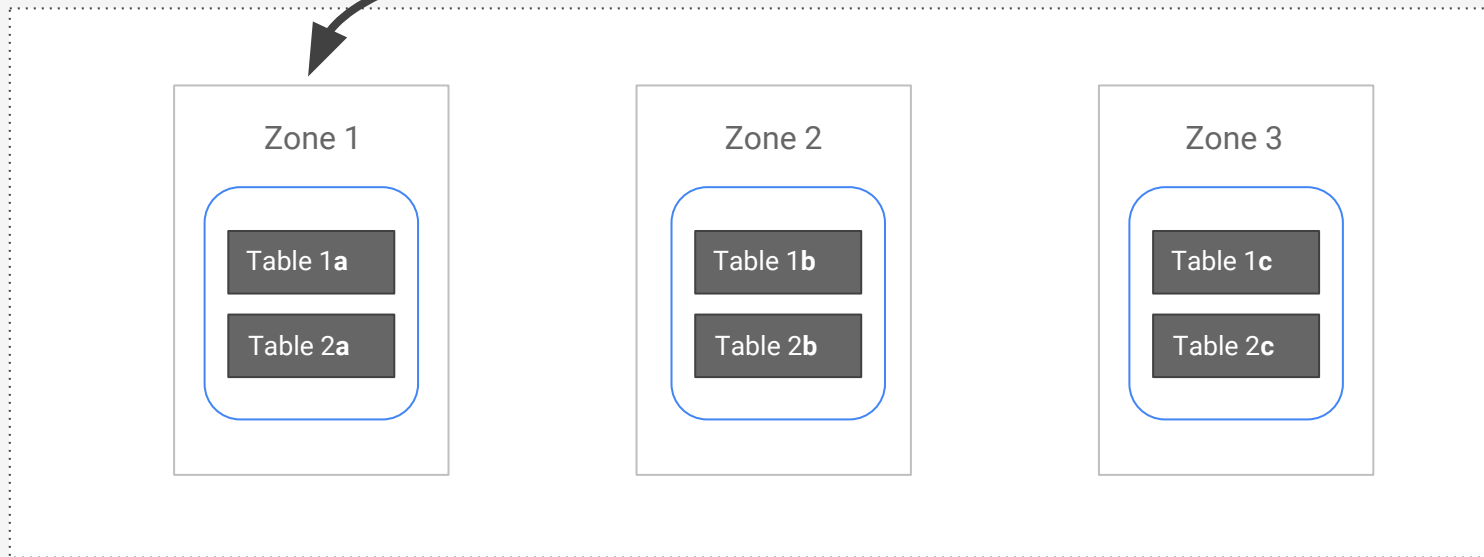
Traditional Database layout -- sharded



Traditional Database layout



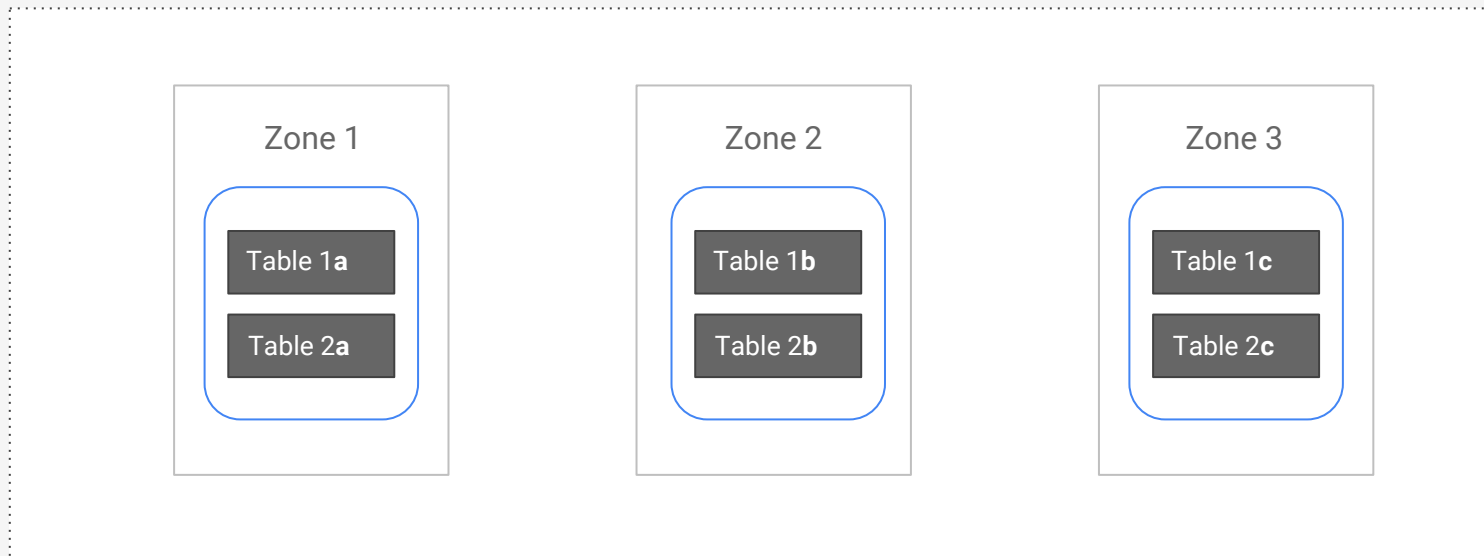
My awesome
app



Traditional Database layout



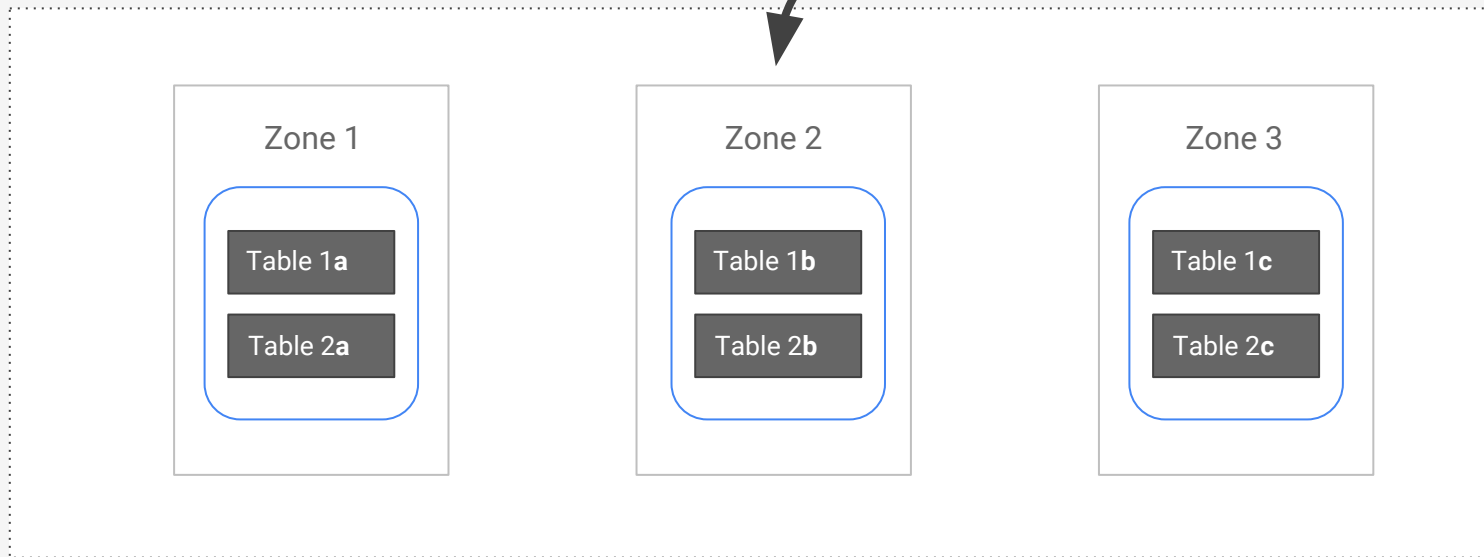
My awesome
app



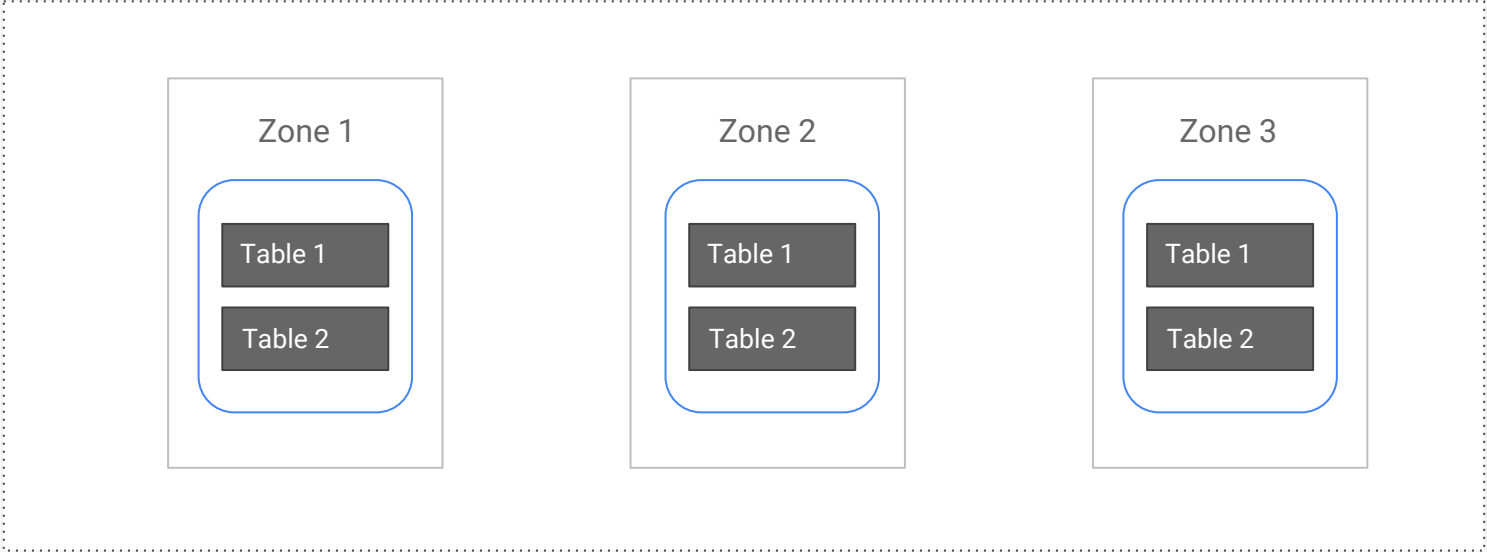
Traditional Database layout



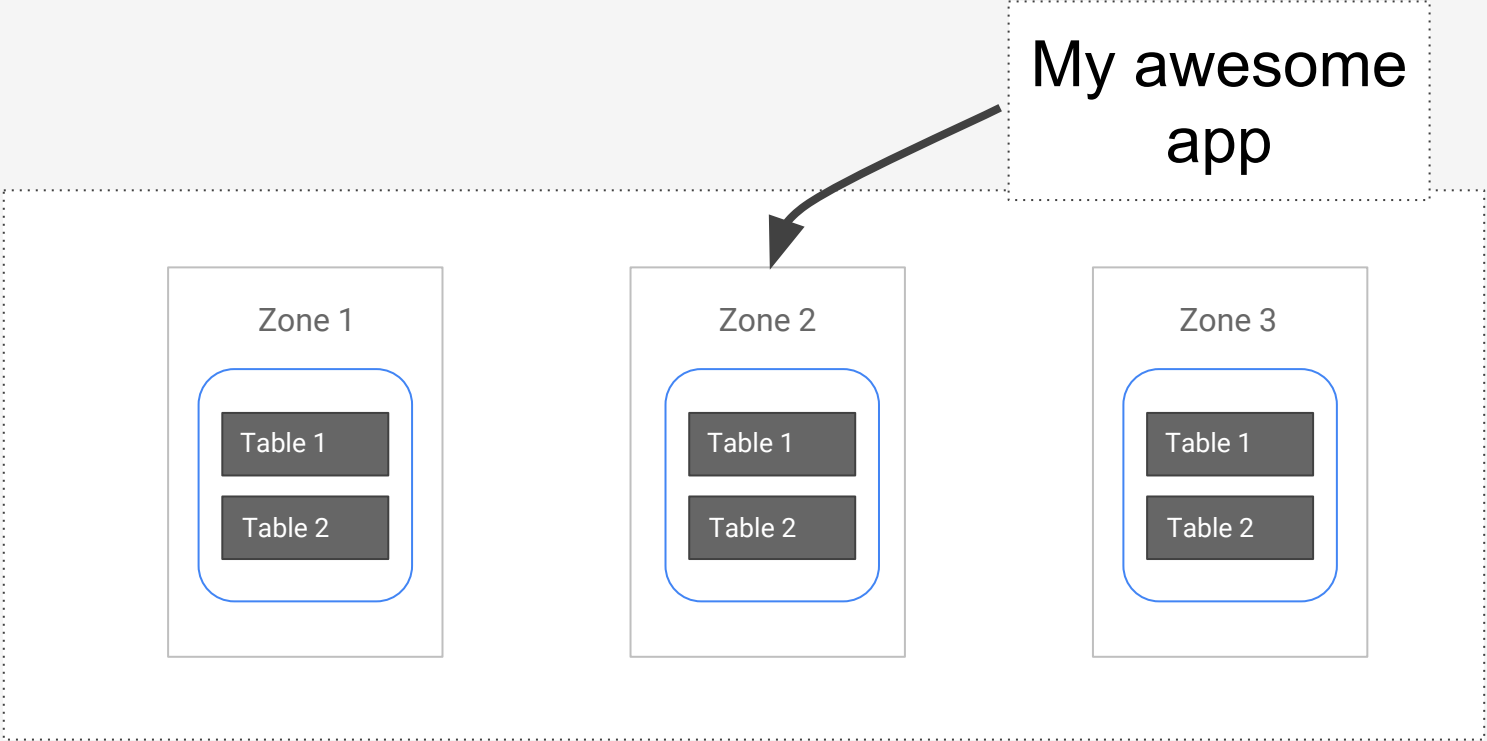
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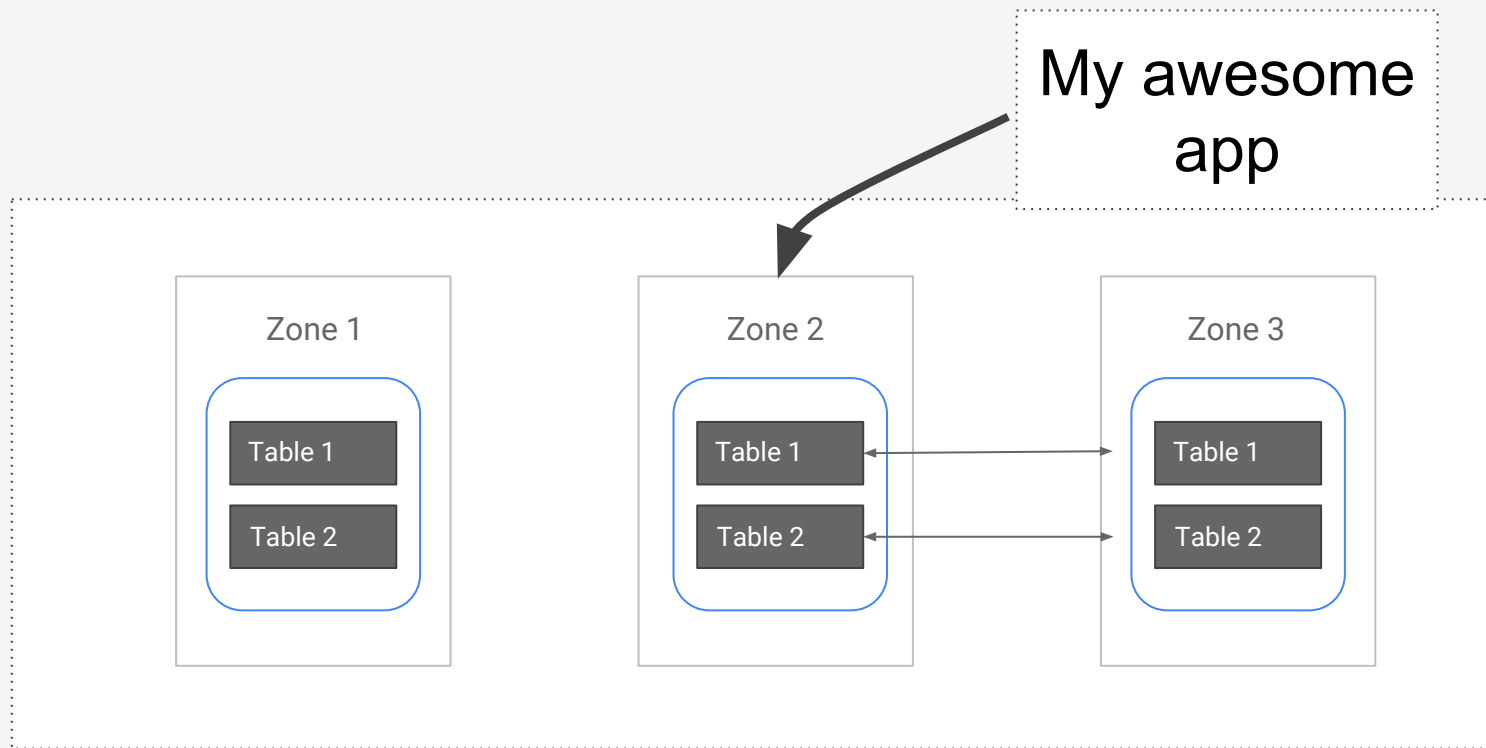
Spanner Database layout



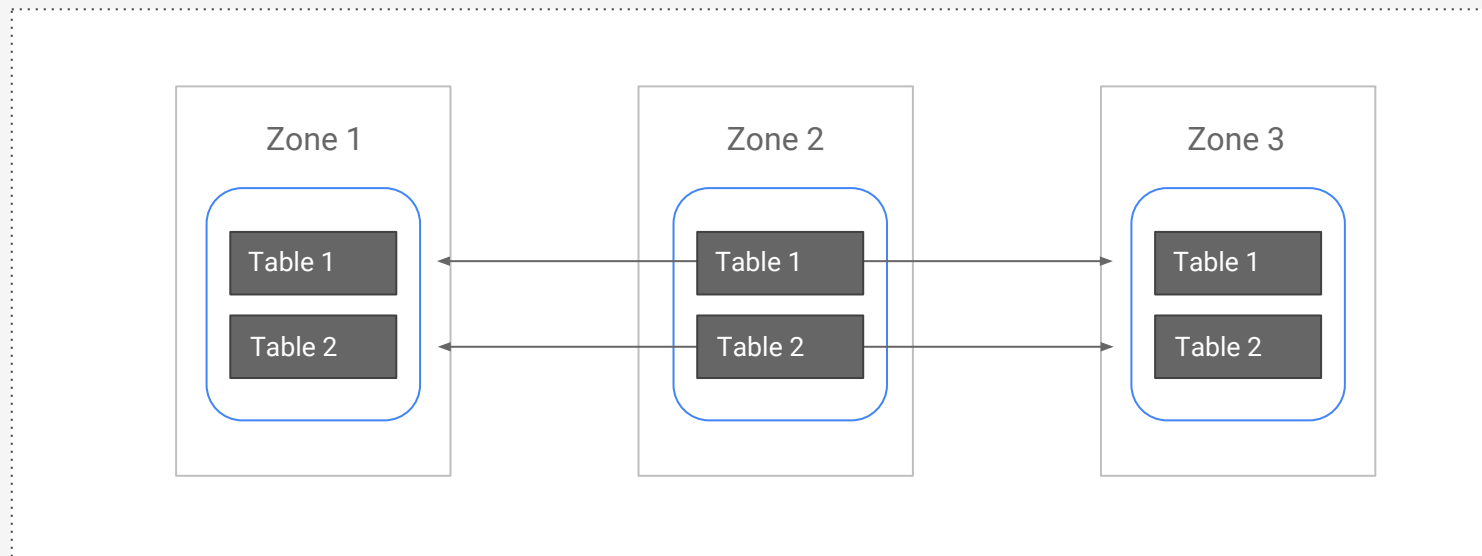
Transactions



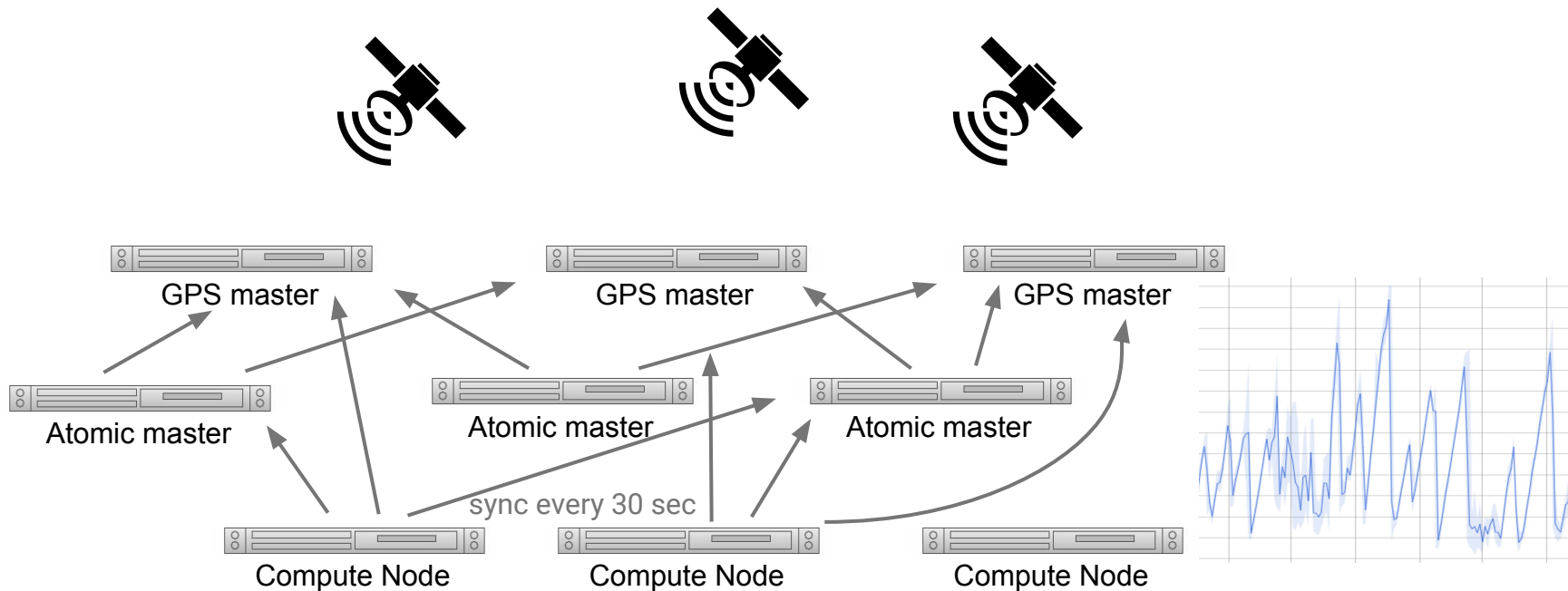
Transactions



Transactions

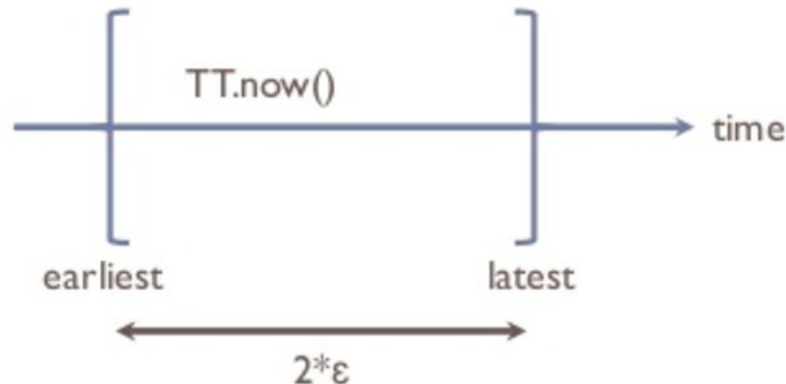


TrueTime Infrastructure

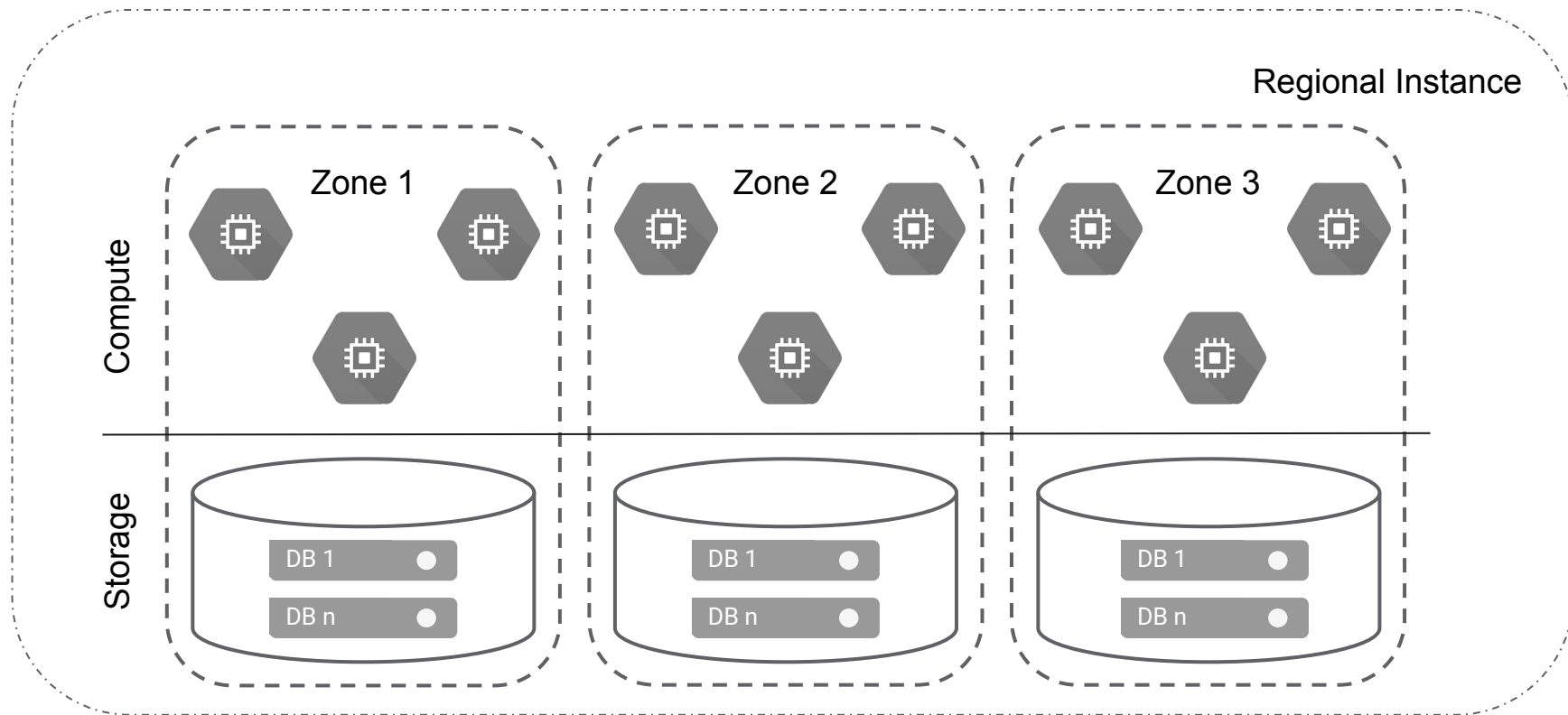


TrueTime - virtual global clock

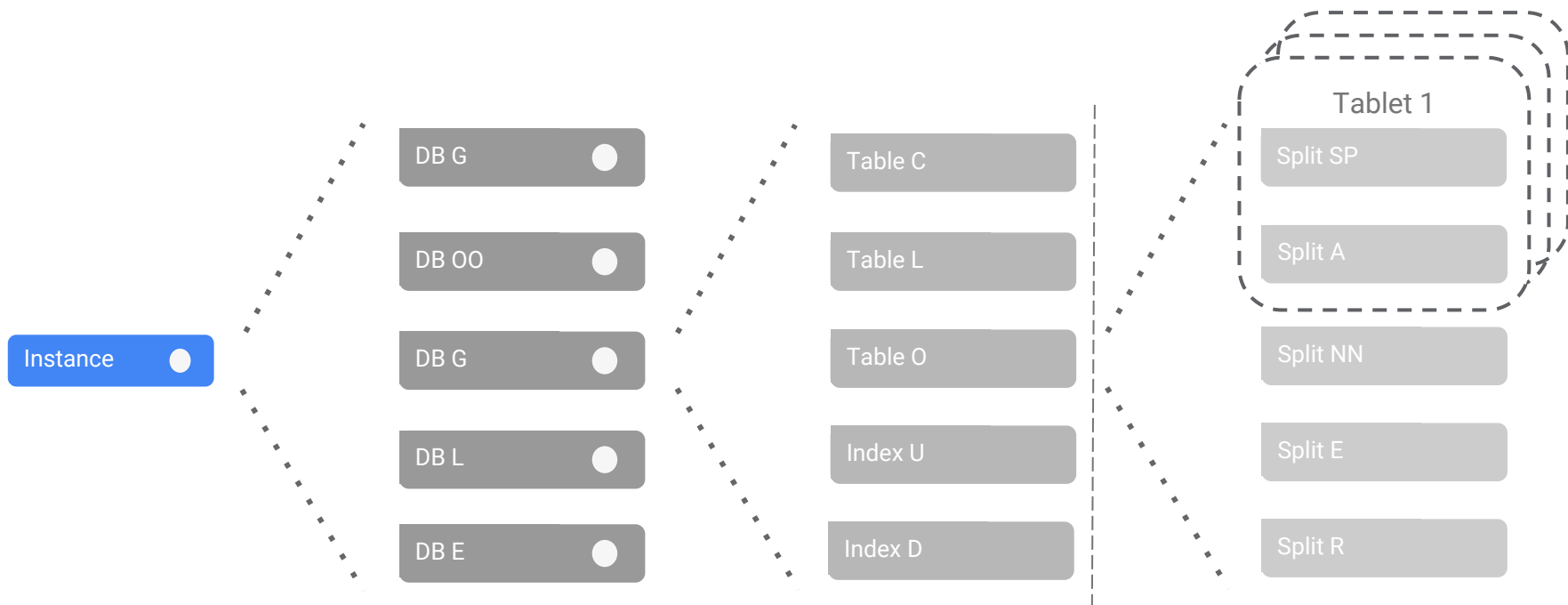
- Quantifies the “worst” possible error / drift between clocks in all datacenters around the world (global clock)
- `TrueTime.Now()` gives you an interval $[t_1, t_2]$; $t_2 = t_1 + 2\epsilon$



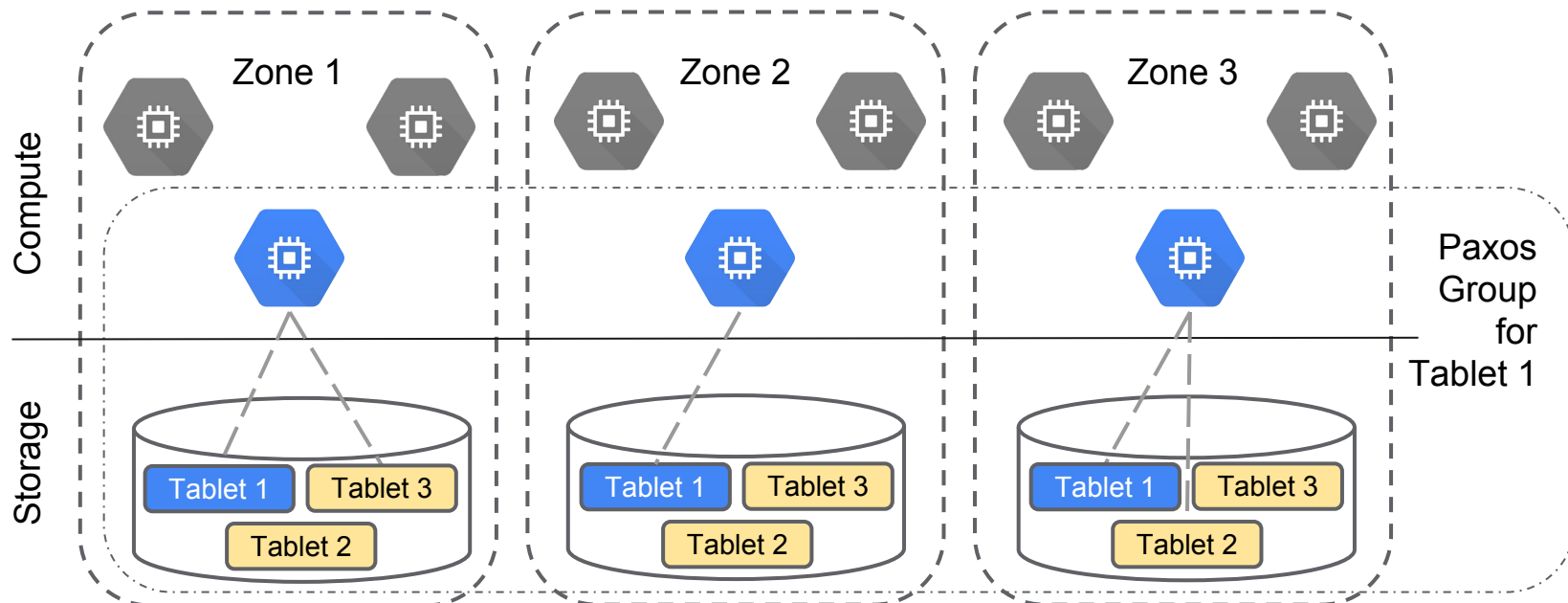
Spanner Internals



Spanner Internals



Spanner Internals

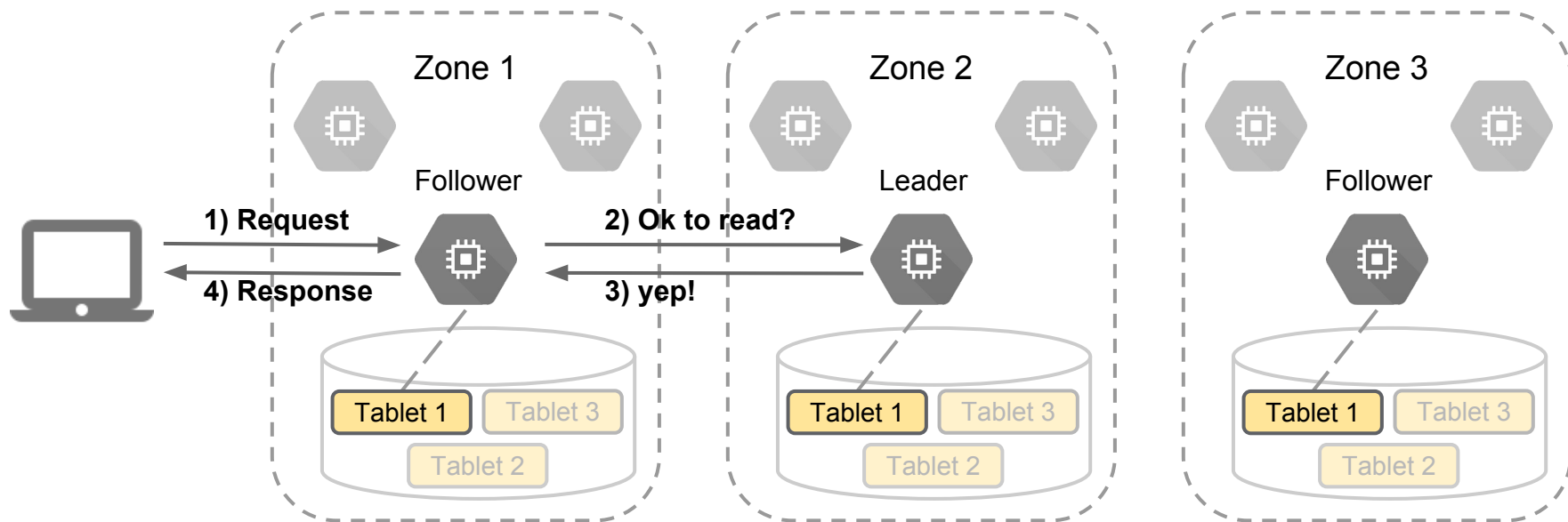


* TrueTime used for leader leases → making sure there is only one leader for a tablet at any given time

Life of a Query

Consistent Read

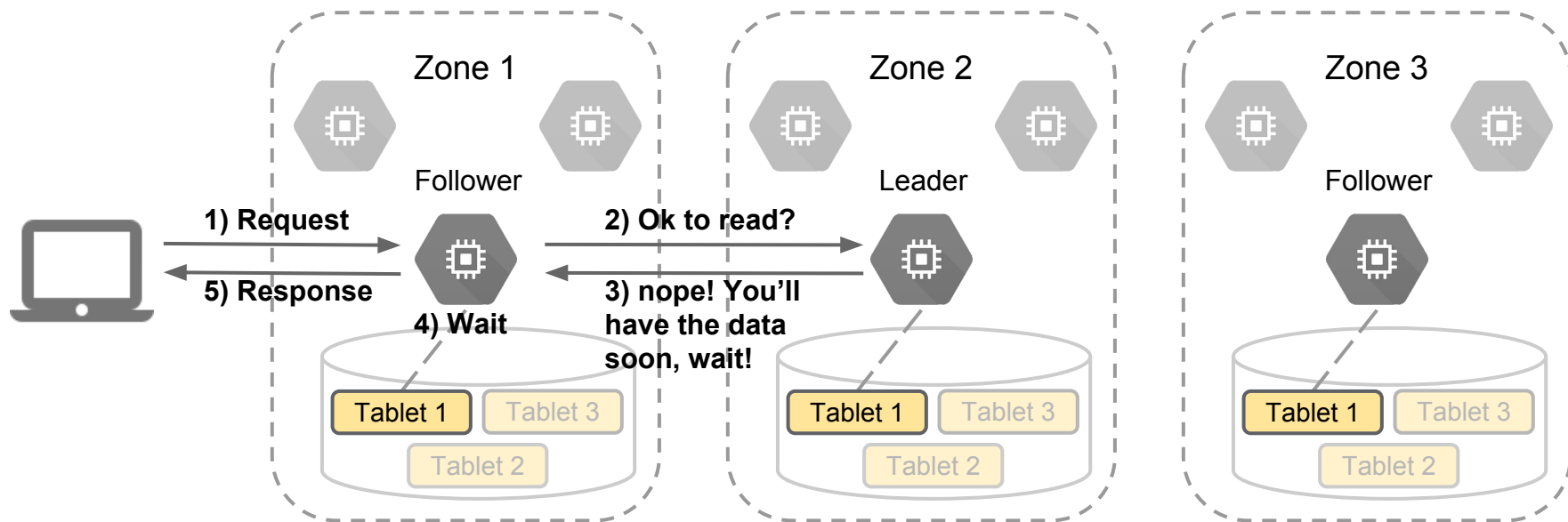
```
SELECT * FROM Company WHERE Name = 'Google';
```



Life of a Query

Consistent Read

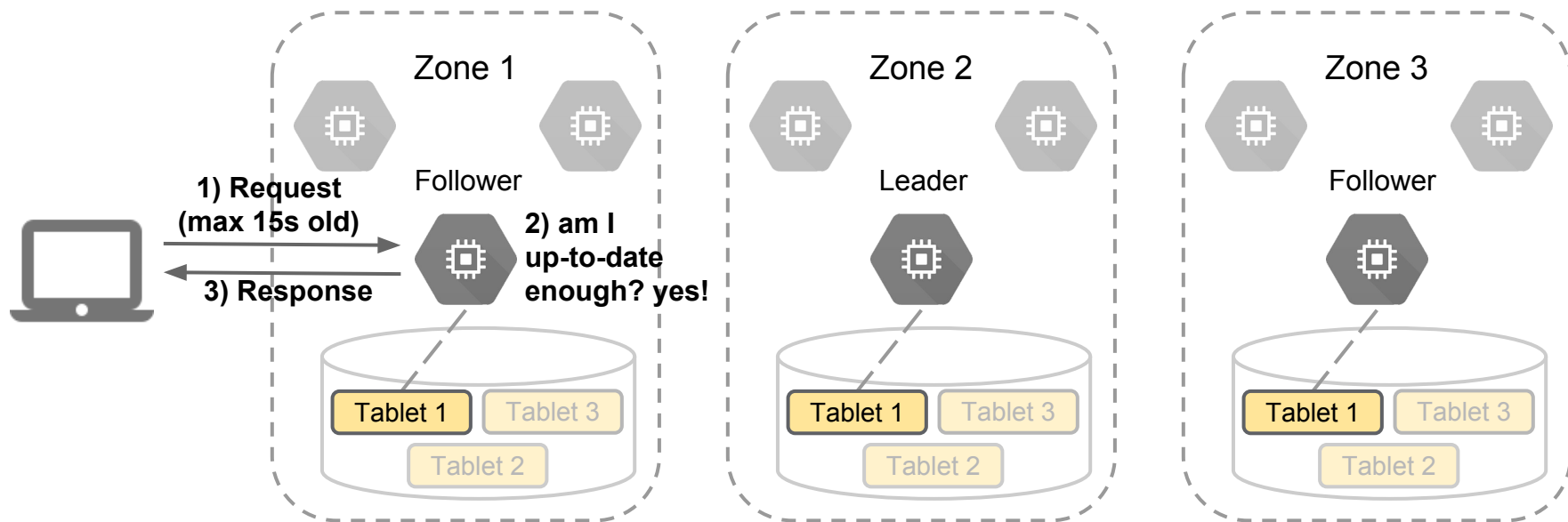
```
SELECT * FROM Company WHERE Name = 'Google';
```



Life of a Query

Time-bounded (stale) reads

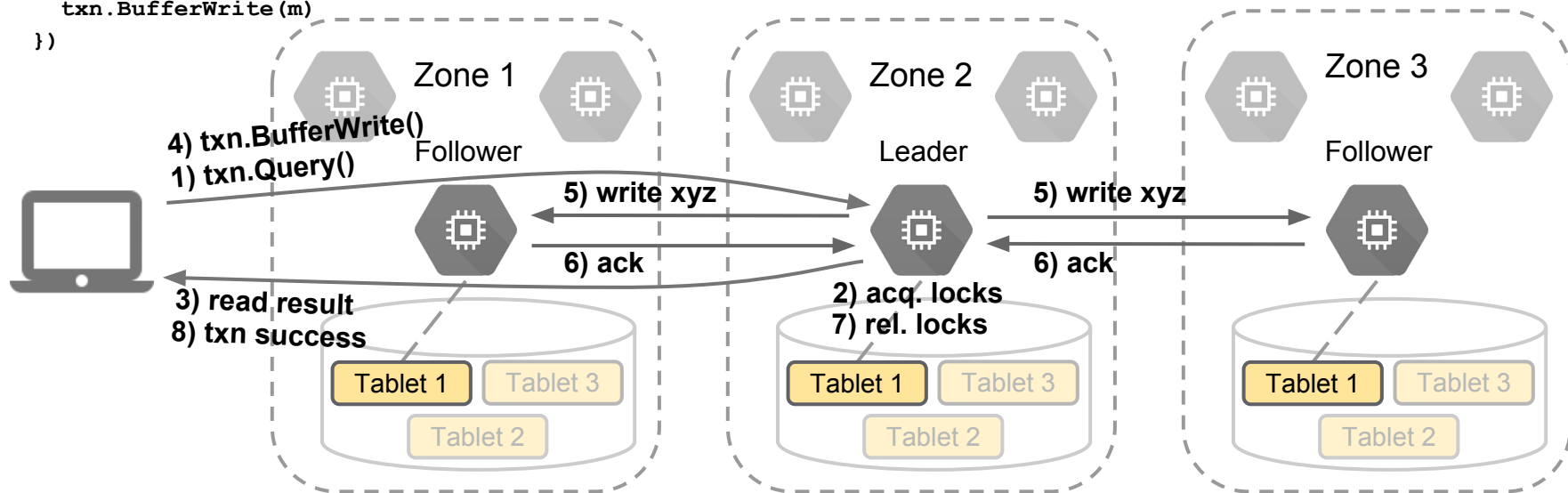
```
SELECT * FROM Company WHERE Name = 'Google';
```



Life of a Query

Read-Write Transaction

```
dbClient.ReadWriteTransaction(ctx, func(txn *spanner.ReadWriteTransaction) error {  
    txn.Query(ctx, stmt).Do(func(r *spanner.Row) error { ... })  
    m = append(m, spanner.InsertOrUpdate("TableName", []string{"ColA", "ColB"},  
        []interface{}{"93bc3d", "Lorem Ipsum"}))  
    txn.BufferWrite(m)  
})
```



Relational Data Layout



SingerId	SingerName
1	Beatles
2	U2
3	Pink Floyd

SingerId	AlbumId	AlbumName
1	1	Help!
1	2	Abbey Road
3	1	The Wall

Interleave Data Layout



1	Beatles	
1	1	Help!
1	2	Abbey Road
2	U2	
3	Pink Floyd	
3	1	The Wall

Relational data model

```
CREATE TABLE Singers (  
  SingerId INT64 NOT NULL,  
  SingerName STRING(MAX),  
) PRIMARY KEY(SingerId);
```

```
CREATE TABLE Albums (  
  SingerId INT64 NOT NULL,  
  AlbumId INT64 NOT NULL,  
  AlbumName STRING(MAX),  
) PRIMARY KEY(SingerId, AlbumId)  
INTERLEAVE IN Singers;
```

No Downtime
Schema Migrations

```
ALTER TABLE Singers  
ADD COLUMN Age INT64;
```

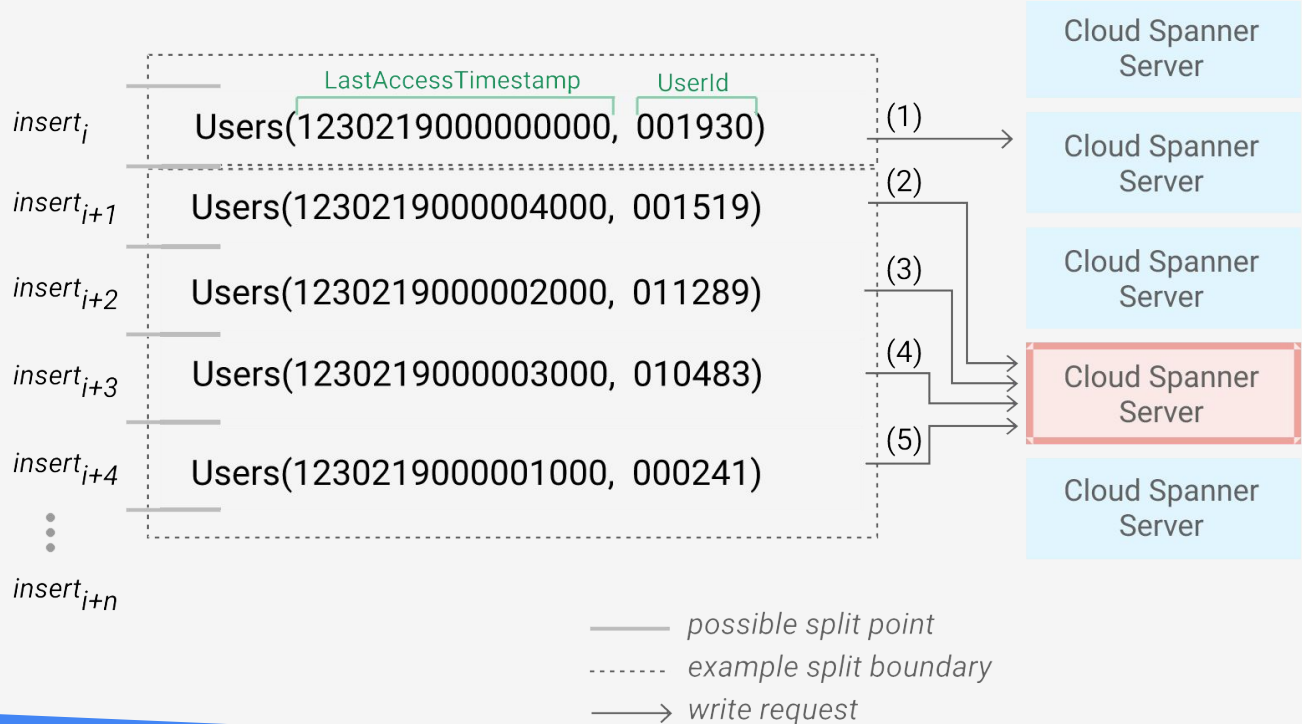
Schema Migration

```
CREATE TABLE `terms` (  
  `id` bigint(19) unsigned NOT NULL AUTO_INCREMENT,  
  `set_id` bigint(19) unsigned NOT NULL,  
  `term` text NOT NULL,  
  `definition` text NOT NULL,  
  `photo` varchar(255) NOT NULL,  
  `term_crc` int(11) unsigned NOT NULL,  
  `rank` int(11) unsigned NOT NULL,  
  `last_modified` int(11) unsigned NOT NULL DEFAULT '0',  
  `is_deleted` tinyint(3) unsigned NOT NULL DEFAULT '0',  
  `term_custom_audio_id` int(11) unsigned DEFAULT NULL,  
  `definition_custom_audio_id` int(11) unsigned DEFAULT NULL,  
  PRIMARY KEY (`id`,`set_id`),  
  KEY `set_id_is_deleted` (`set_id`,`is_deleted`)  
) ENGINE=InnoDB AUTO_INCREMENT=0 DEFAULT CHARSET=utf8mb4 ROW_FORMAT=COMPRESSED  
/*150100 PARTITION BY HASH (set_id) PARTITIONS 25 */
```

```
CREATE TABLE terms (  
  id INT64 NOT NULL,  
  set_id INT64 NOT NULL,  
  spanner_id STRING(36) NOT NULL,      -- v4 uuid  
  spanner_set_id STRING(36) NOT NULL,  -- v4 uuid  
  term STRING(MAX) NOT NULL,  
  definition STRING(MAX) NOT NULL,  
  image STRING(1024) NOT NULL,  
  term_crc INT64 NOT NULL,  
  rank INT64 NOT NULL,  
  last_modified TIMESTAMP NOT NULL,    -- defaults to 0  
  is_deleted INT64 NOT NULL,           -- defaults to 0  
  term_custom_audio_id INT64,  
  definition_custom_audio_id INT64  
) PRIMARY KEY(spanner_set_id, spanner_id);  
  
CREATE INDEX terms_id on terms(id);
```

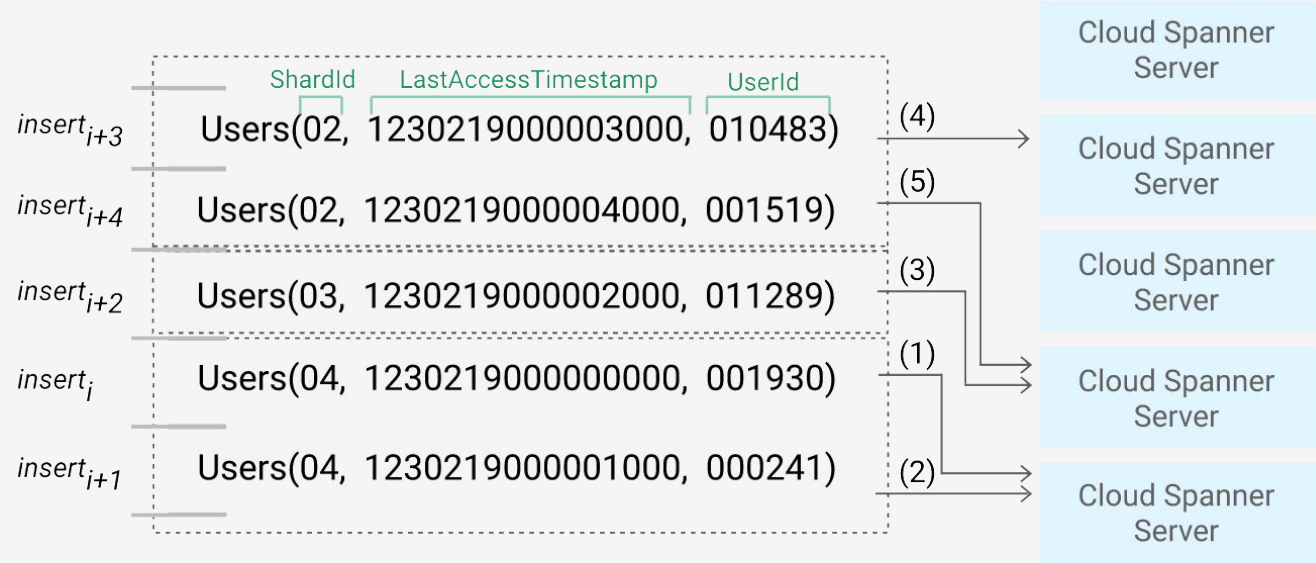


Anti-Pattern: Timestamp Ordering



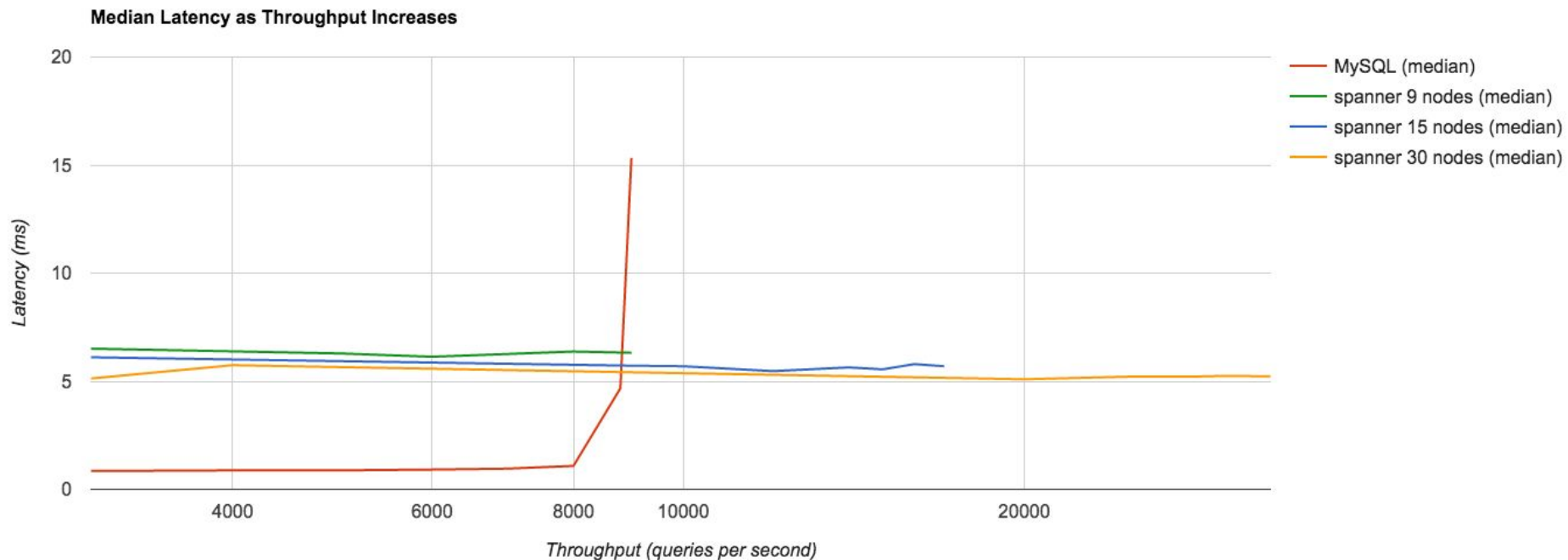


Anti-Pattern: Sequences



— possible split point
- - - - - example split boundary
—> write request

Median Latency





Demo

Thank You

@joe_int